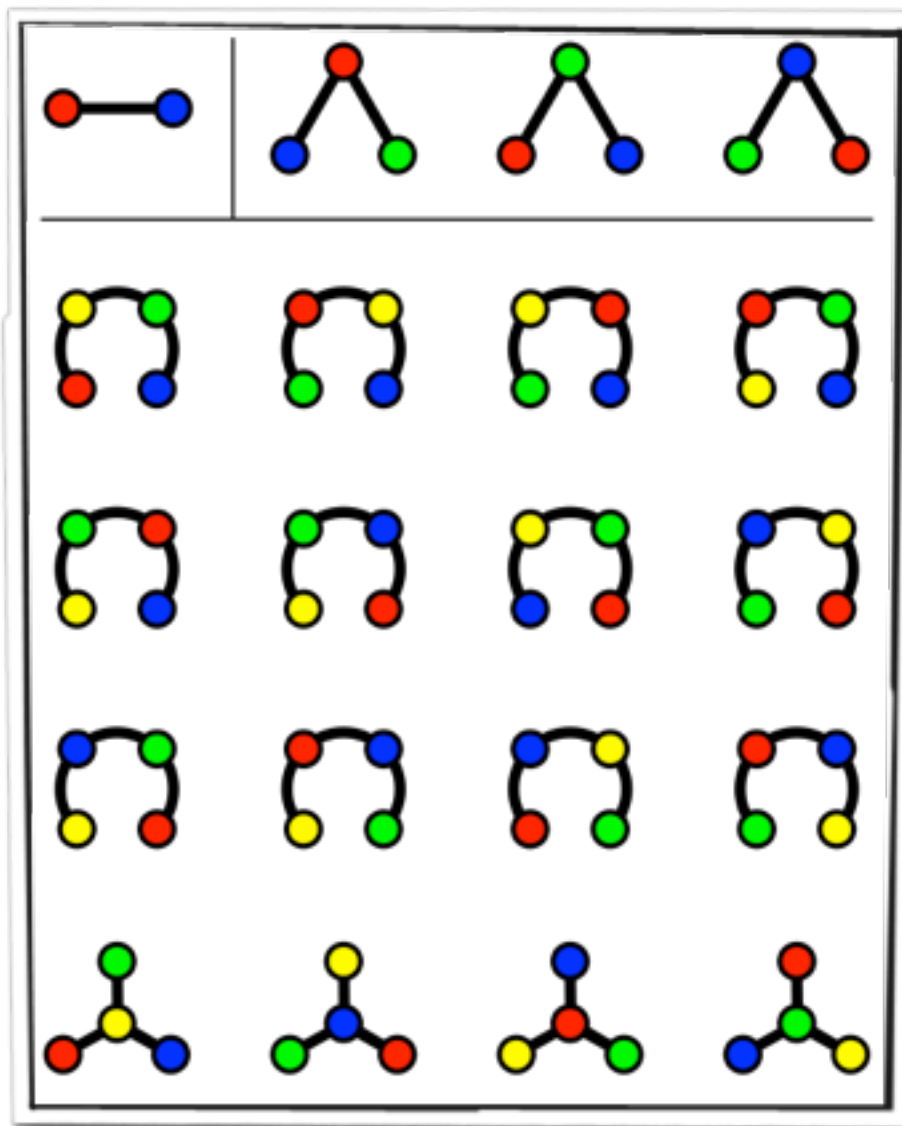


# Combinatorics

南京大学  
尹一通

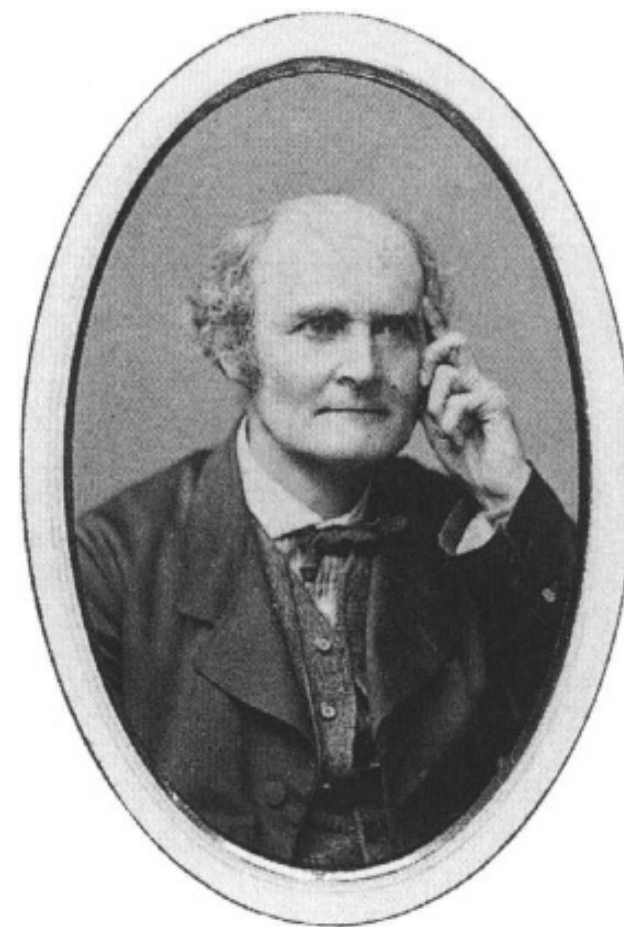
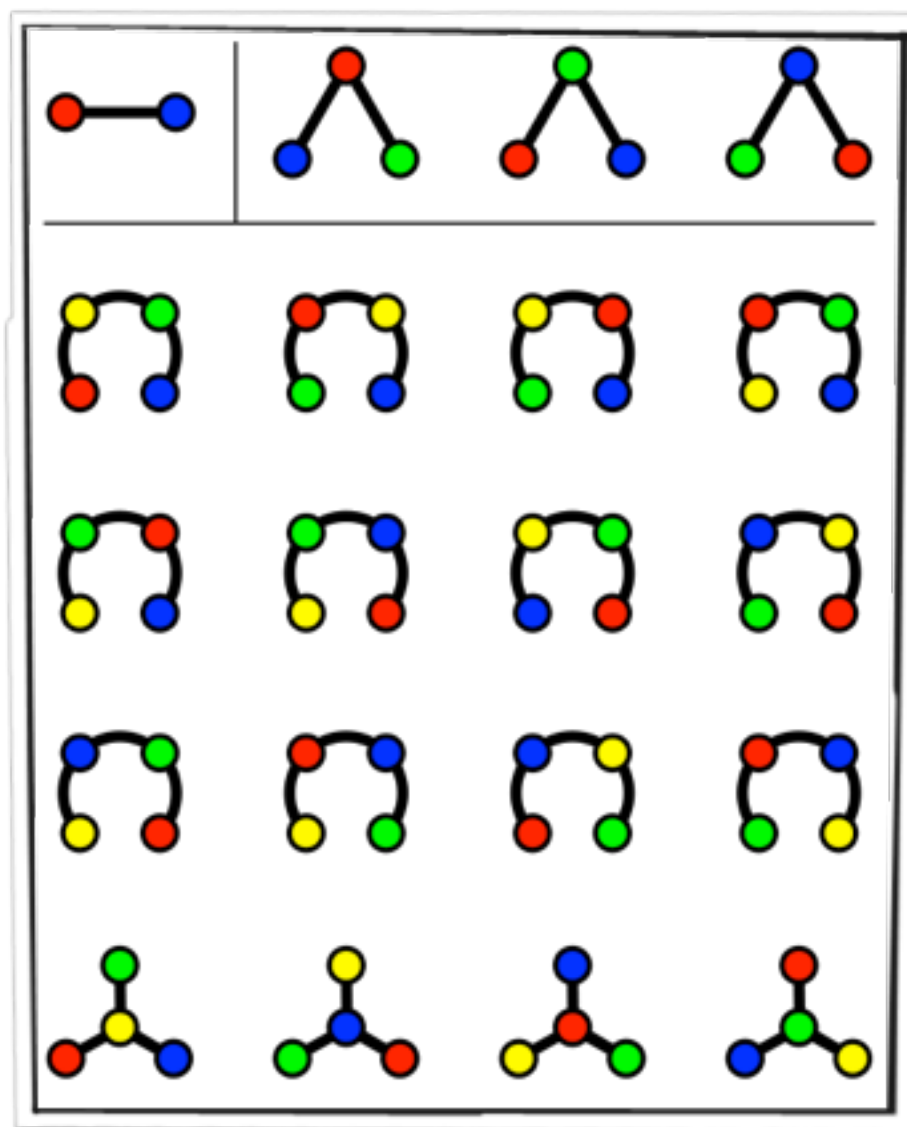
# Counting (labeled) trees



*“How many different trees  
can be formed from  
 $n$  **distinct** vertices?”*

# Cayley's formula:

There are  $n^{n-2}$  trees on  $n$  distinct vertices.



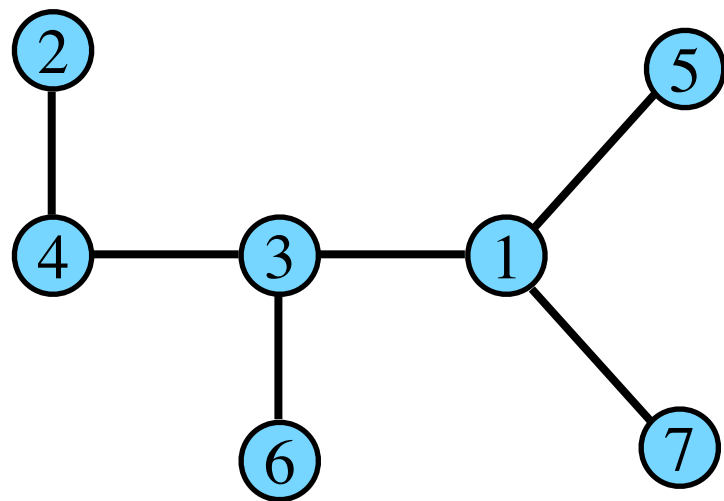
Arthur Cayley

# Prüfer Code

leaf : vertex of degree 1

removing a leaf from  $T$ , still a tree

$T_1$ :



$u_i$ : 2, 4, 5, 6, 3, 1

$v_i$ : 4, 3, 1, 3, 1, 7

$T_1 = T$ ;

for  $i = 1$  to  $n-1$

$u_i$ : smallest leaf in  $T_i$ ;

$(u_i, v_i)$ : edge in  $T_i$ ;

$T_{i+1} = \text{delete } u_i \text{ from } T_i$ ;

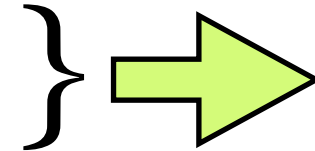
Prüfer code:

$(v_1, v_2, \dots, v_{n-2})$

edges of  $T : (u_i, v_i), 1 \leq i \leq n-1$

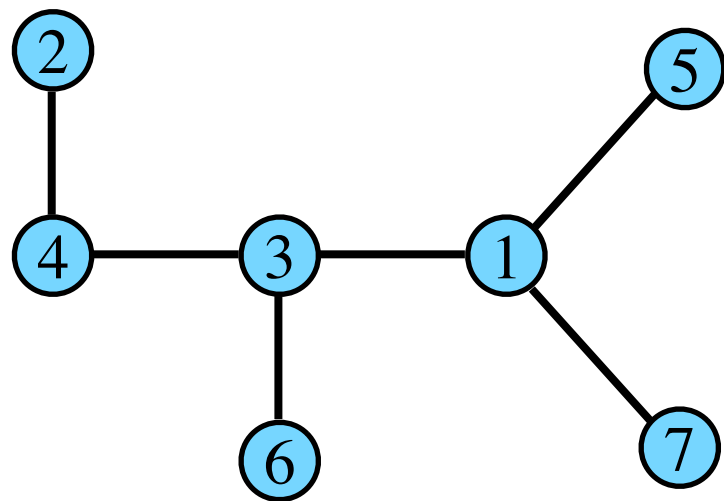
$$v_{n-1} = n$$

$u_i$ : smallest leaf in  $T_i$   
a tree has  $\geq 2$  leaves



$n$  is never deleted  
 $u_i \neq n$

$T :$



$u_i : 2, 4, 5, 6, 3, 1$

$v_i : \underline{4, 3, 1, 3, 1, 7}$

$(v_1, v_2, \dots, v_{n-2})$

Only need to recover  
every  $u_i$  from  $(v_1, v_1, \dots, v_{n-2})$ .

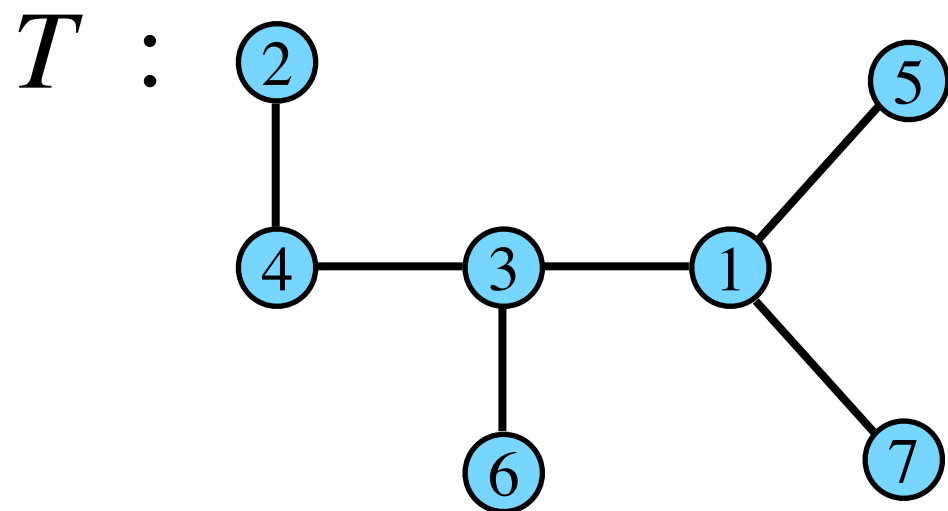
$u_i$  is the smallest number not in  
 $\{u_1, \dots, u_{i-1}\} \cup \{v_i, \dots, v_{n-1}\}$

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 $\{u_1, \dots, u_{i-1}\} \cup \{v_i, \dots, v_{n-1}\}$

$\forall$  vertex  $v$  in  $T$ ,

# occurrences of  $v$  in  $u_1, u_2, \dots, u_{n-1}, v_{n-1}$  : **1**

# occurrences of  $v$  in edges  $(u_i, v_i), 1 \leq i \leq n-1$ :  **$\deg_T(v)$**



# occurrences of  $v$  in  
**Prüfer code**:  $(v_1, v_2, \dots, v_{n-2})$

**$\deg_T(v)-1$**

$u_i$ : 2, 4, 5, 6, 3, 1

$v_i$ : 4, 3, 1, 3, 1, 7

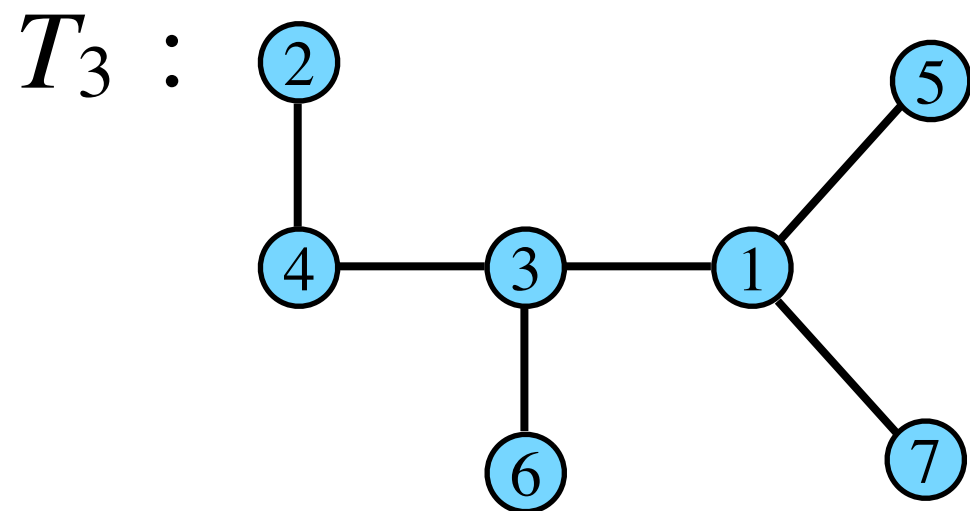
$(v_1, v_2, \dots, v_{n-2})$

$u_i$  is the smallest number not in  
 $\{u_1, \dots, u_{i-1}\} \cup \{v_i, \dots, v_{n-1}\}$

$\forall$  vertex  $v$  in  $T_i$ ,

# occurrences of  $v$  in  $u_i, u_{i+1}, \dots, u_{n-1}, v_{n-1}$  : 1

# occurrences of  $v$  in edges  $(u_j, v_j), i \leq j \leq n-1$ :  $\deg_{T_i}(v)$



# occurrences of  $v$  in  $(v_i, \dots, v_{n-2})$

$\deg_{T_i}(v) - 1$

leaf  $v$  of  $T_i$ :

in  $\{u_i, u_{i+1}, \dots, u_{n-1}, v_{n-1}\}$

not in  $\{v_i, v_{i+1}, \dots, v_{n-2}\}$

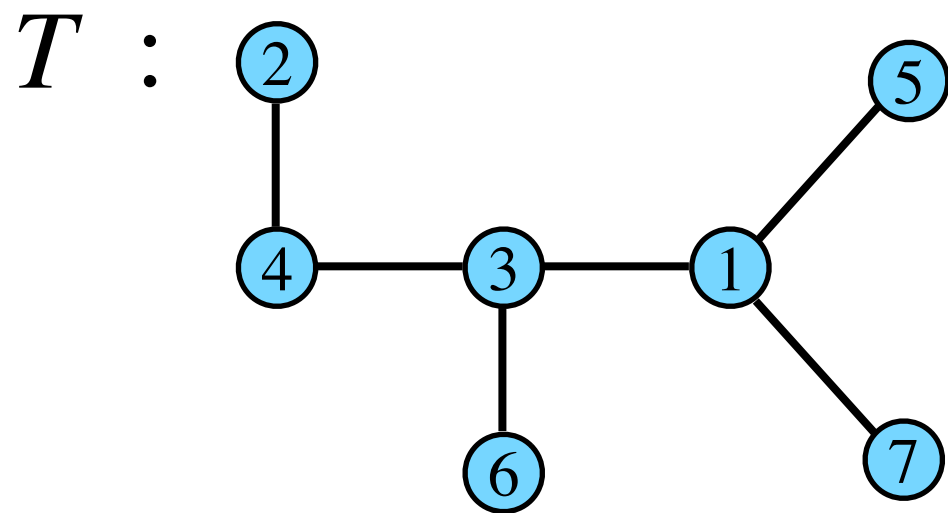
$u_i$ : 2, 4, 5, 6, 3, 1

$v_i$ : 4, 3, 1, 3, 1, 7

$(v_1, v_2, \dots, v_{n-2})$

$u_i$ : smallest leaf in  $T_i$

$u_i$  is the smallest number not in  
 $\{u_1, \dots, u_{i-1}\} \cup \{v_i, \dots, v_{n-1}\}$



$u_i$ : 2, 4, 5, 6, 3, 1

$v_i$ : 4, 3, 1, 3, 1, 7

$(v_1, v_2, \dots, v_{n-2})$

$T =$  empty graph;

$v_{n-1} = n$  ;

for  $i = 1$  to  $n-1$

$u_i$ : smallest number not in  
 $\{u_1, \dots, u_{i-1}\} \cup \{v_i, \dots, v_{n-1}\}$

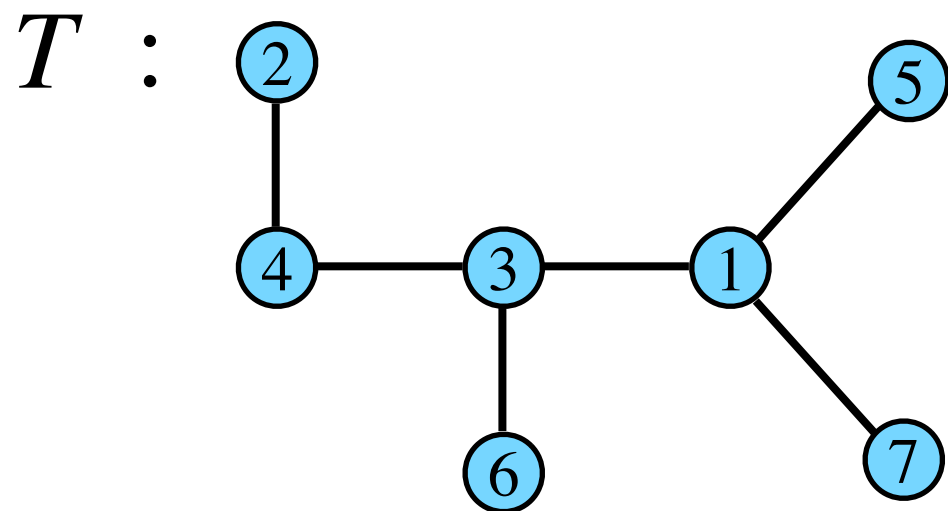
add edge  $(u_i, v_i)$  to  $T$ ;



Prüfer code is reversible  $\Rightarrow$  1-1

every  $(v_1, v_2, \dots, v_{n-2}) \in \{1, 2, \dots, n\}^{n-2}$

is decodable to a tree  $\Rightarrow$  onto



$u_i$ : 2, 4, 5, 6, 3, 1

$v_i$ : 4, 3, 1, 3, 1, 7

$(v_1, v_2, \dots, v_{n-2})$

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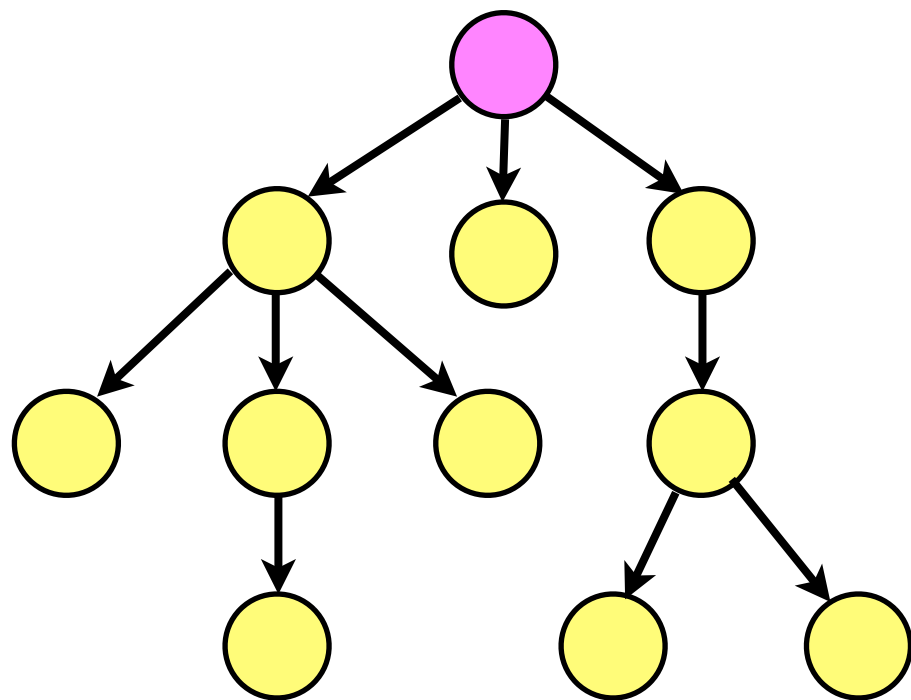
is decodable to a tree  onto

### **Cayley's formula:**

There are  $n^{n-2}$  trees on  $n$  distinct vertices.

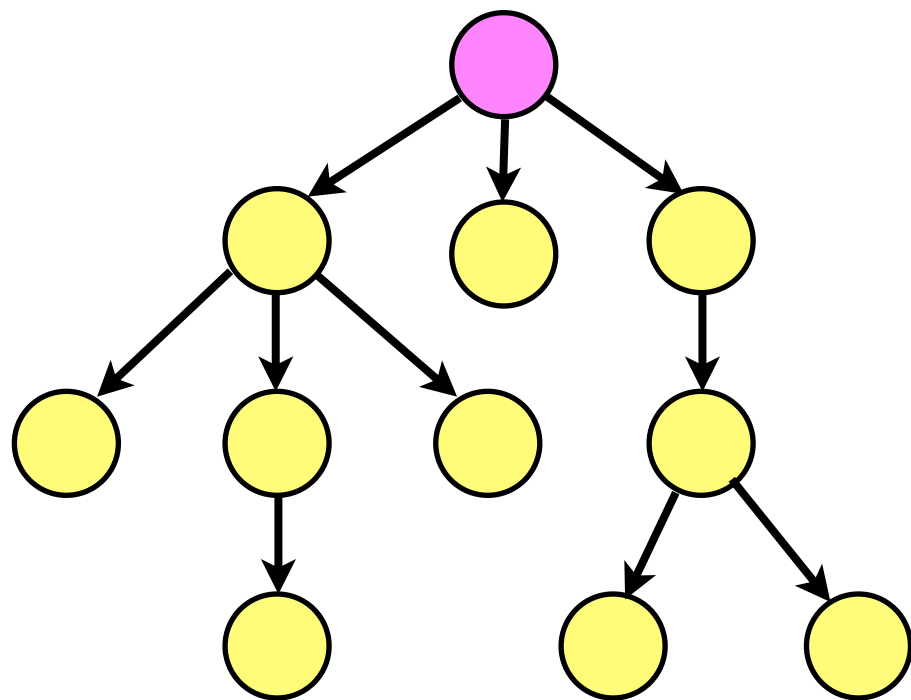
# Double Counting

# of *sequences* of adding *directed edges* to an empty graph to form a *rooted tree*



$T_n$  : # of trees on  $n$  distinct vertices.

# of *sequences* of adding *directed edges* to an empty graph to form a *rooted tree*



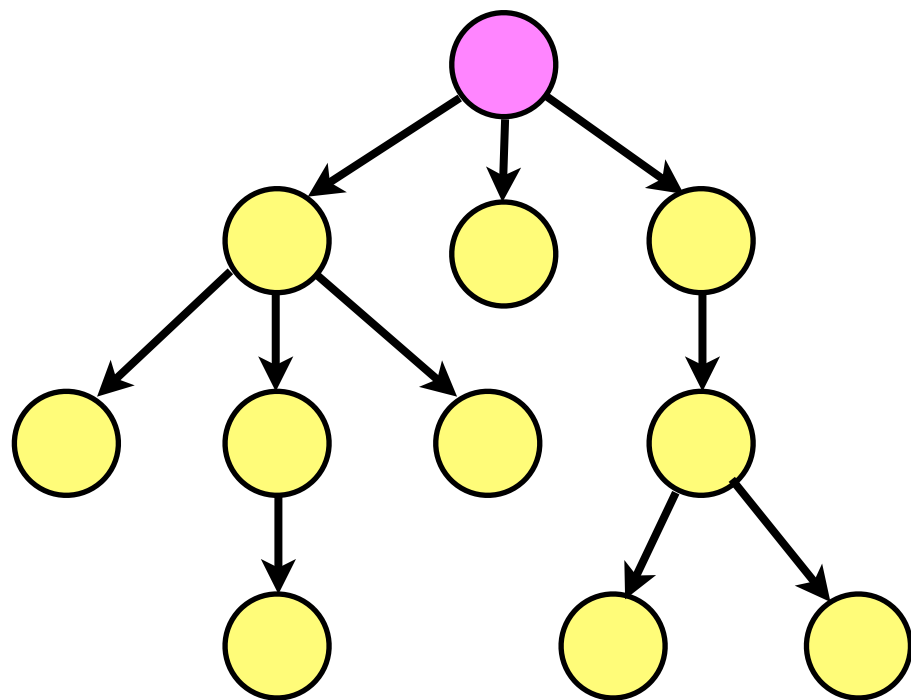
From a tree:

- pick a root;
- pick an order of edges.

$$T_n n(n-1)!$$
$$= n! T_n$$

$T_n$  : # of trees on  $n$  distinct vertices.

# of *sequences* of adding *directed edges* to an empty graph to form a *rooted tree*

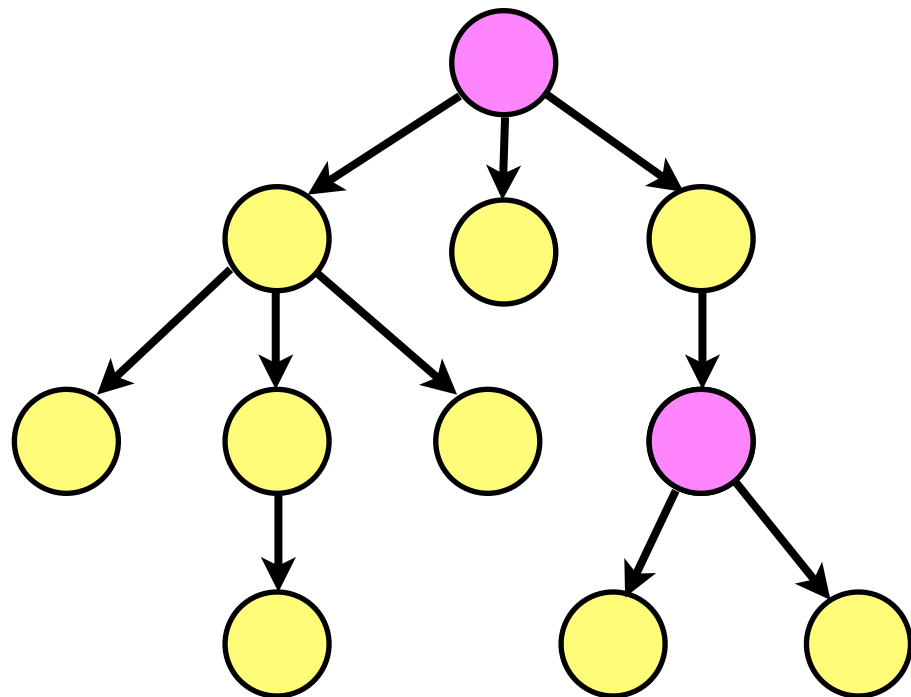


From an empty graph:

- add edges one by one

# of *sequences* of adding *directed edges* to an empty graph to form a *rooted tree*

From an empty graph: • add edges one by one

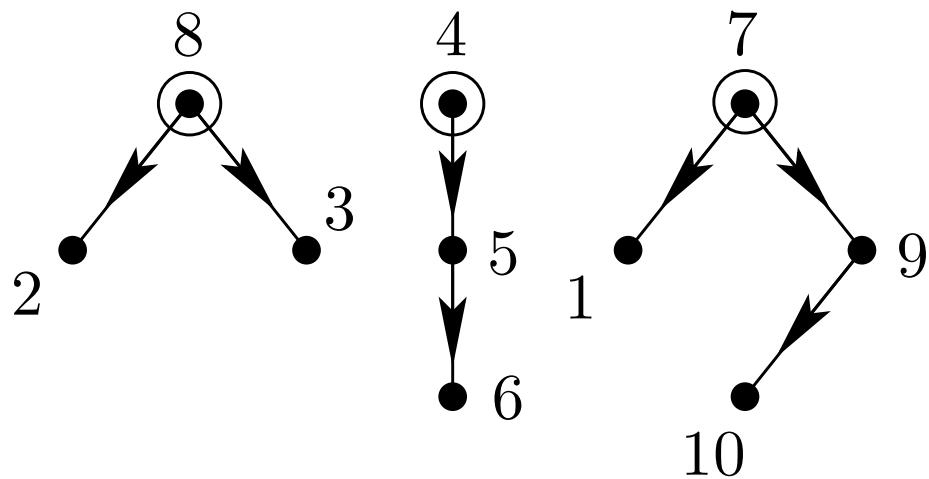


Start from  $n$  ~~isolated vertices~~  
*rooted trees*

Each step joins 2 trees.

# of *sequences* of adding *directed edges* to an empty graph to form a *rooted tree*

From an empty graph: • add edges one by one

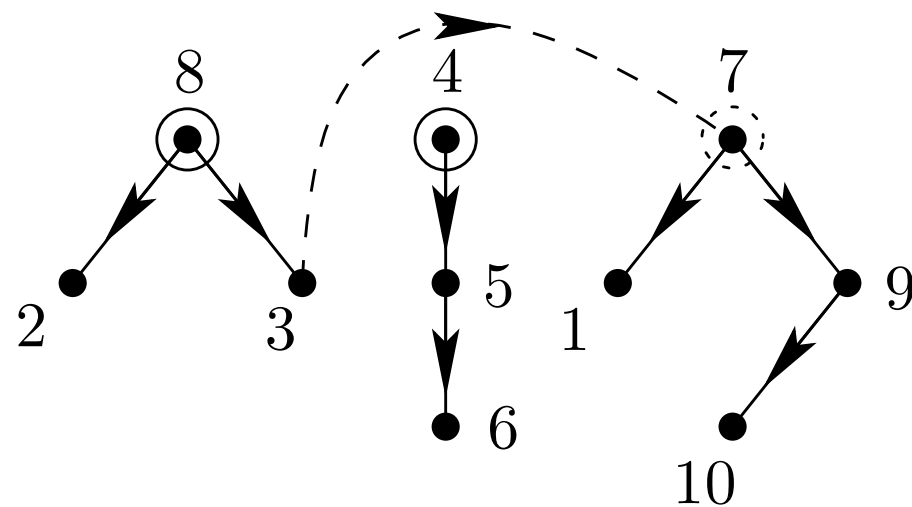


Start from  $n$  rooted trees.

After adding  $k$  edges

$n-k$  rooted trees

add an edge



any  
vertex

$n$



root of  
another tree

$n-k-1$

# of *sequences* of adding *directed edges* to an empty graph to form a *rooted tree*

From an empty graph: • add edges one by one

$$\prod_{k=0}^{n-2} n(n-k-1)$$

Start from  $n$  rooted trees.

After adding  $k$  edges

$n-k$  rooted trees

add an edge

$$= n^{n-1} \prod_{k=1}^{n-1} k$$

$$= n^{n-2} n!$$

any  
vertex

$n$



root of  
another tree

$n-k-1$



# of *sequences* of adding *directed edges* to an empty graph to form a *rooted tree*

From a tree:

- pick a root;
- pick an order of edges.

$$T_n n(n-1)!$$
$$= n! T_n$$

=

From an empty graph:

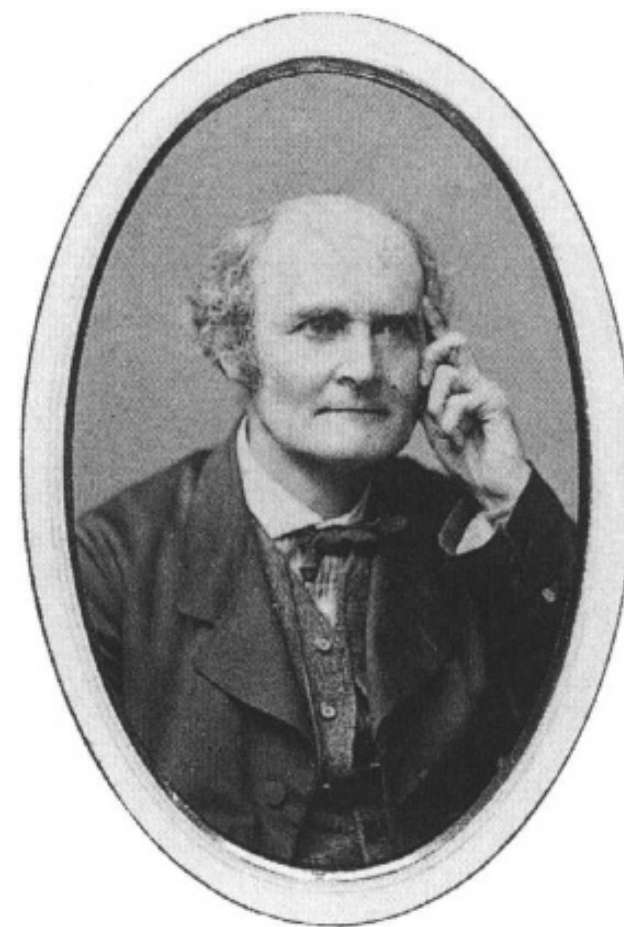
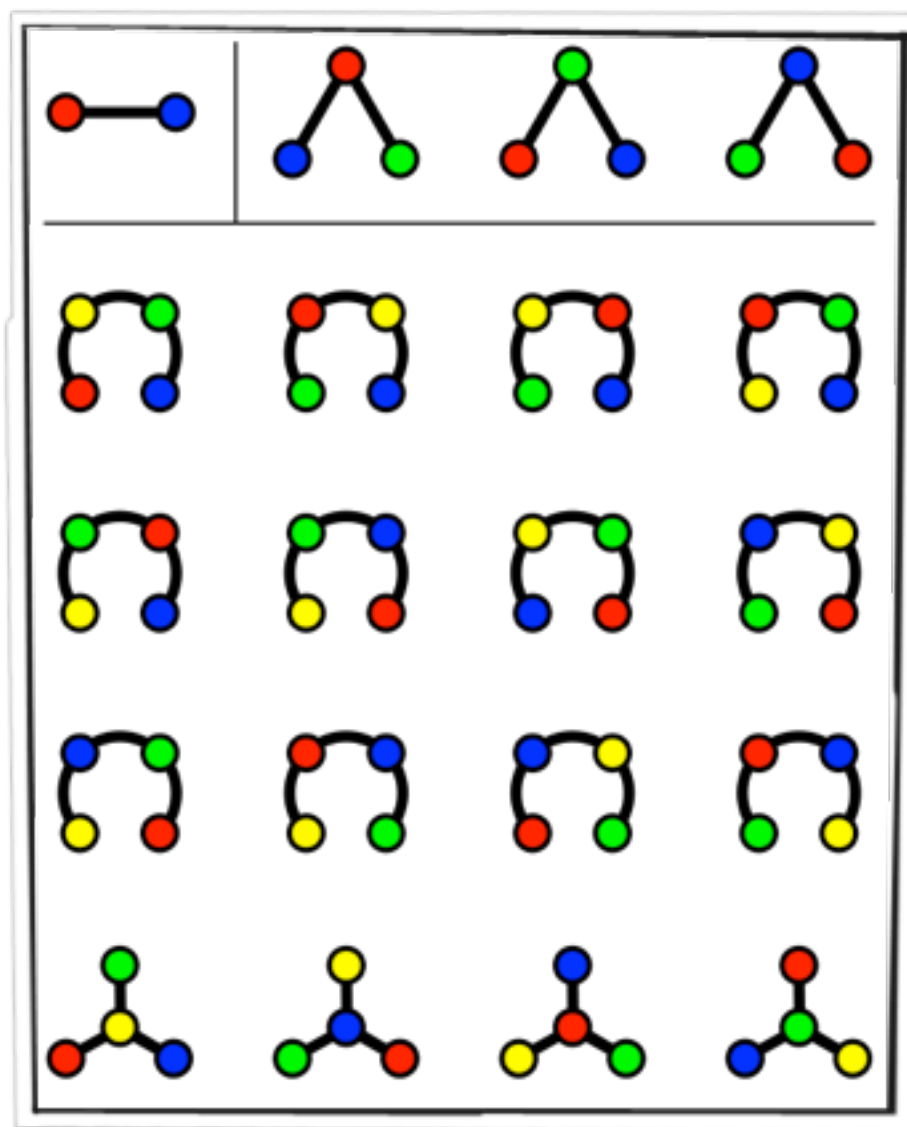
- add edges one by one

$$\prod_{k=2}^n n(k-1)$$
$$= n^{n-2} n!$$

$$T_n = n^{n-2}$$

# Cayley's formula:

There are  $n^{n-2}$  trees on  $n$  distinct vertices.



Arthur Cayley

# Graph *Laplacian*

Graph  $G(V,E)$

adjacency matrix  $A$

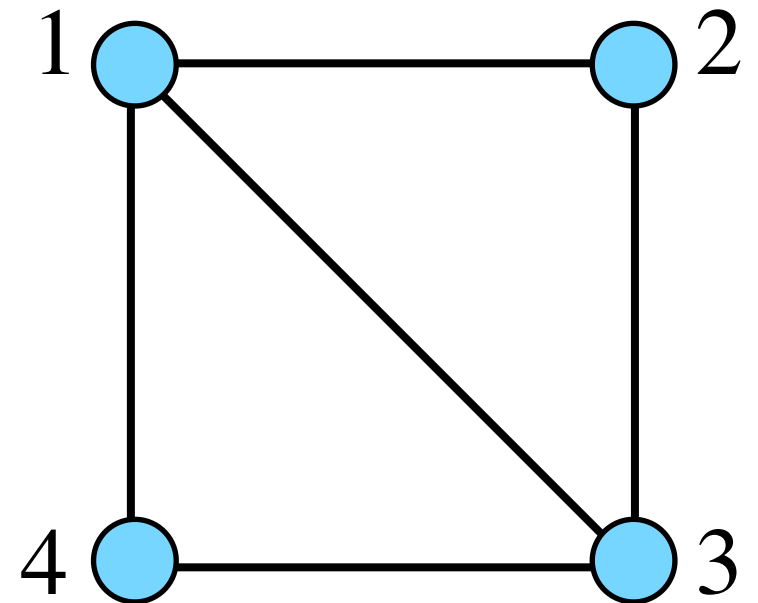
$$A(i, j) = \begin{cases} 1 & \{i, j\} \in E \\ 0 & \{i, j\} \notin E \end{cases}$$

diagonal matrix  $D$

$$D(i, j) = \begin{cases} \deg(i) & i = j \\ 0 & i \neq j \end{cases}$$

**graph *Laplacian*  $L$**

$$L = D - A$$



$$D = \begin{bmatrix} d_1 & & & \\ & d_2 & & \\ & & 0 & \\ & 0 & \ddots & \\ & & & d_n \end{bmatrix}$$

$$L = \begin{bmatrix} 3 & -1 & -1 & -1 \\ -1 & 2 & -1 & 0 \\ -1 & -1 & 3 & -1 \\ -1 & 0 & -1 & 2 \end{bmatrix}$$

# Graph *Laplacian*

graph *Laplacian*  $L$

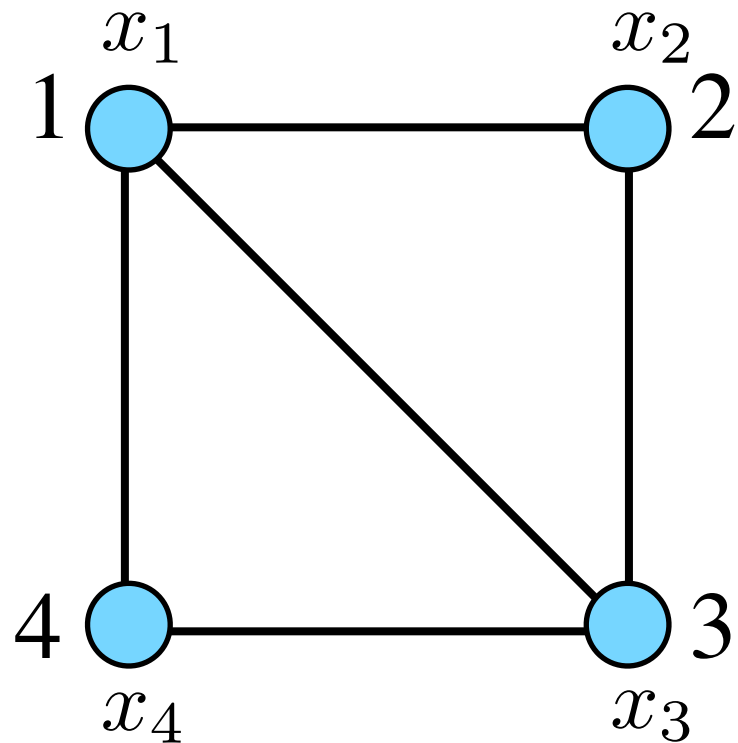
$$L(i, j) = \begin{cases} \deg(i) & i = j \\ -1 & i \neq j, \{i, j\} \in E \\ 0 & \text{otherwise} \end{cases}$$

quadratic form:

$$xLx^T = \sum_i d_i x_i^2 - \sum_{ij \in E} x_i x_j = \frac{1}{2} \sum_{ij \in E} (x_i - x_j)^2$$

incidence matrix  $B : n \times m$

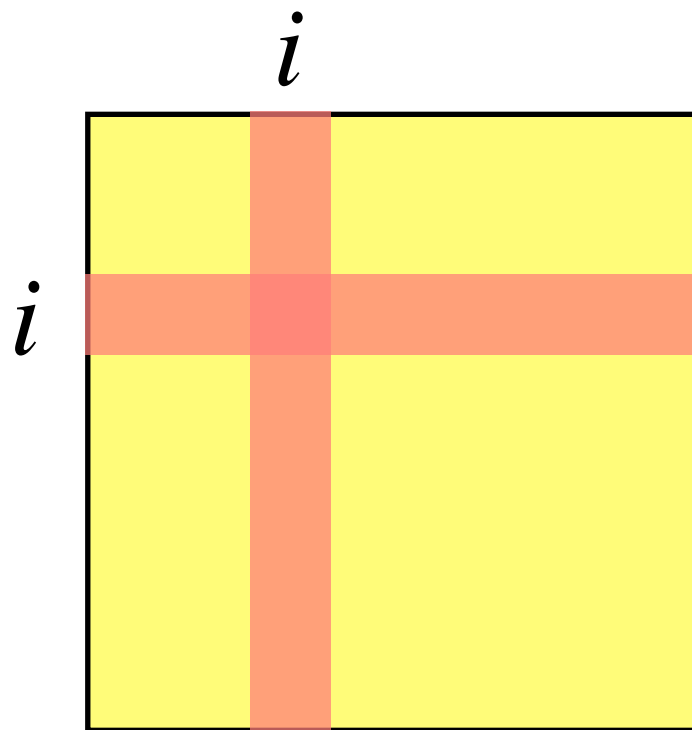
$$i \in V, e \in E \quad B(i, e) = \begin{cases} 1 & e = \{i, j\}, i < j \\ -1 & e = \{i, j\}, i > j \\ 0 & \text{otherwise} \end{cases}$$



$$L = BB^T$$

# *Kirchhoff's* matrix-tree theorem

$L_{i,i}$  : submatrix of  $L$  by removing  
 $i$ th row and  $i$ th column



$t(G)$  : number of spanning trees in  $G$

# *Kirchhoff's* matrix-tree theorem

$L_{i,i}$  : submatrix of  $L$  by removing  
ith row and ith column

$t(G)$  : number of spanning trees in  $G$

**Kirchhoff's Matrix-Tree Theorem:**

$$\forall i, \quad t(G) = \det(L_{i,i})$$

## **Kirchhoff's Matrix-Tree Theorem:**

$$\forall i, \quad t(G) = \det(L_{i,i})$$

$$B_i : (n - 1) \times m$$

incidence matrix  $B$  removing  $i$ th row

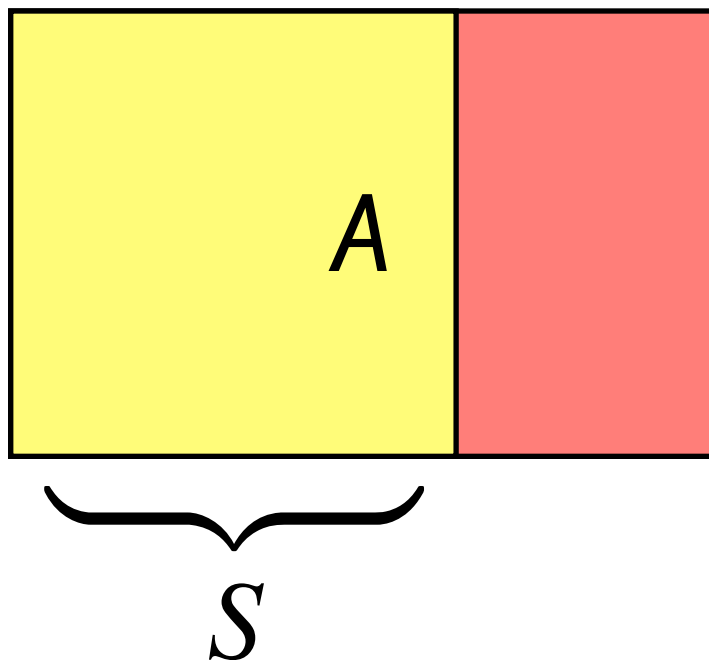
$$L = BB^T$$

$$L_{i,i} = B_i B_i^T \quad \det(L_{i,i}) = \det(B_i B_i^T) = ?$$

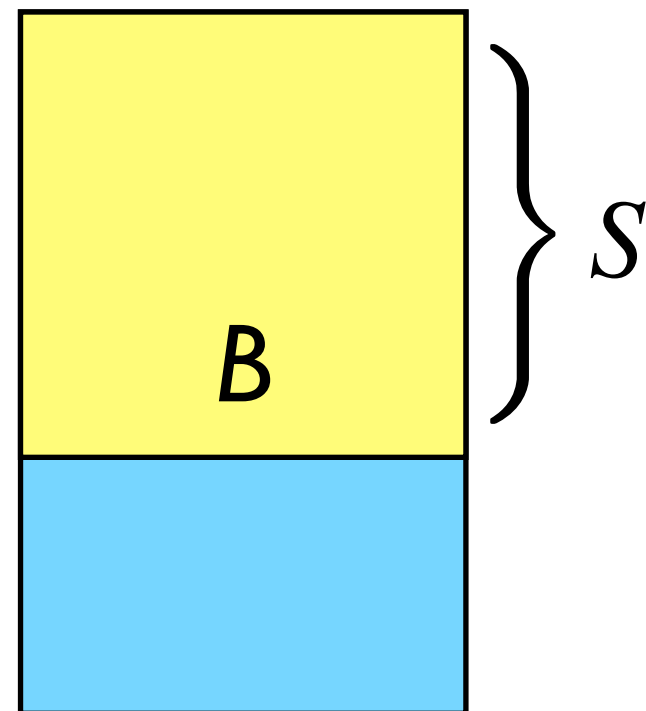
## Cauchy-Binet Theorem:

$$\det(AB) = \sum_{S \in \binom{[m]}{n}} \det(A_{[n], S}) \det(B_{S, [n]})$$

$A : n \times m$



$B : m \times n$





## Cauchy-Binet Theorem:

$$\det(AB) = \sum_{S \in \binom{[m]}{n}} \det(A_{[n],S}) \det(B_{S,[n]})$$

$$\det(L_{i,i}) = \det(B_i B_i^T)$$

$$= \sum_{S \in \binom{[m]}{n-1}} \det(B_{[n] \setminus \{i\}, S}) \det(B_{S, [n] \setminus \{i\}}^T)$$

$$= \sum_{S \in \binom{[m]}{n-1}} \det(B_{[n] \setminus \{i\}, S})^2$$

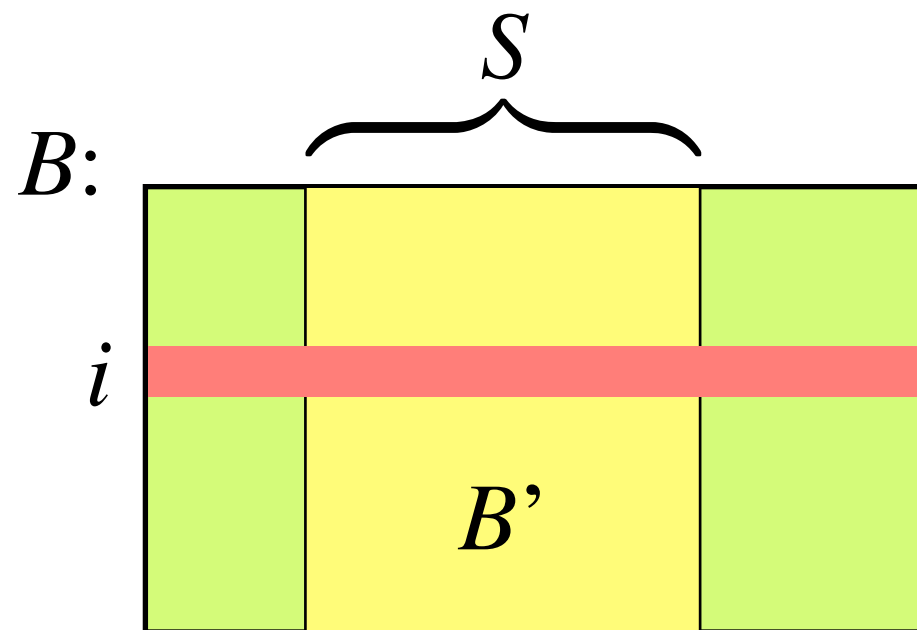
$$\det(L_{i,i}) = \sum_{S \in \binom{[m]}{n-1}} \det(B_{[n] \setminus \{i\}, S})^2$$

$$j \in [n] \setminus \{i\}, e \in S$$

$$B_{[n] \setminus \{i\}, S}(j, e) = \begin{cases} 1 & e = \{j, k\}, j < k \\ -1 & e = \{j, k\}, j > k \\ 0 & \text{otherwise} \end{cases}$$

$$\det(B_{[n] \setminus \{i\}, S}) = \begin{cases} \pm 1 & S \text{ is a spanning tree of } G \\ 0 & \text{otherwise} \end{cases}$$

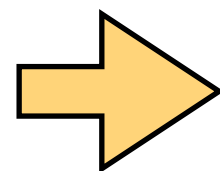
$$\det(B_{[n]\setminus\{i\},S}) = \begin{cases} \pm 1 & S \text{ is a spanning tree of } G \\ 0 & \text{otherwise} \end{cases}$$



$$B' = B_{[n]\setminus\{i\},S}$$

$(n-1) \times (n-1)$  **matrix:**

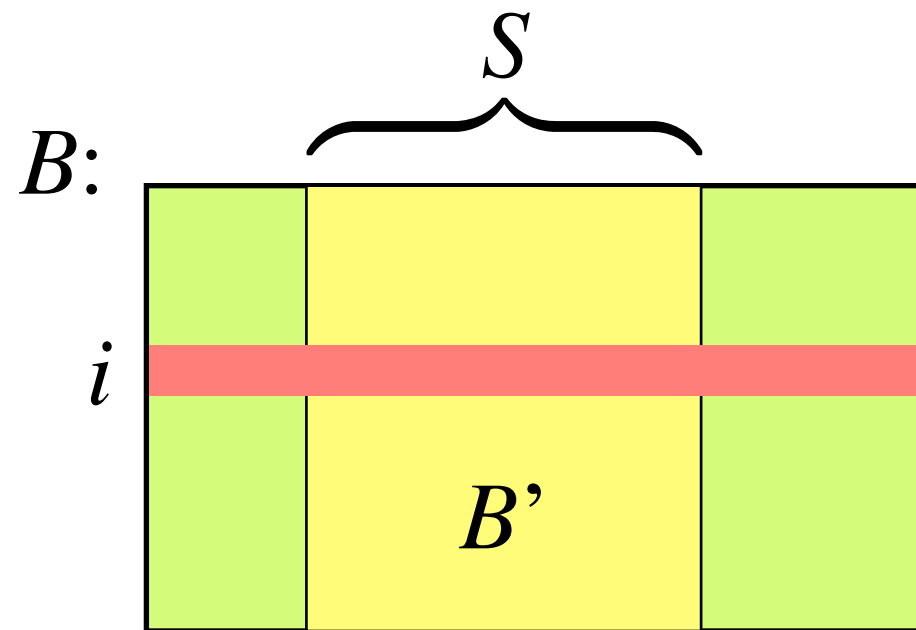
every column contains  
at most one 1 and at most one -1  
and all other entries are 0



$$\det(B') \in \{-1, 0, 1\}$$

$\det(B') \neq 0$  iff  $S$  is a spanning tree

$\det(B') \neq 0$  iff  $S$  is a spanning tree



$S$  is not a spanning tree:

$\exists$  a connected component  $R$   
s.t.  $i \notin R$

$\Rightarrow \det(B') = 0$

$S$  is a spanning tree:

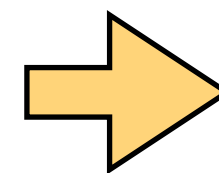
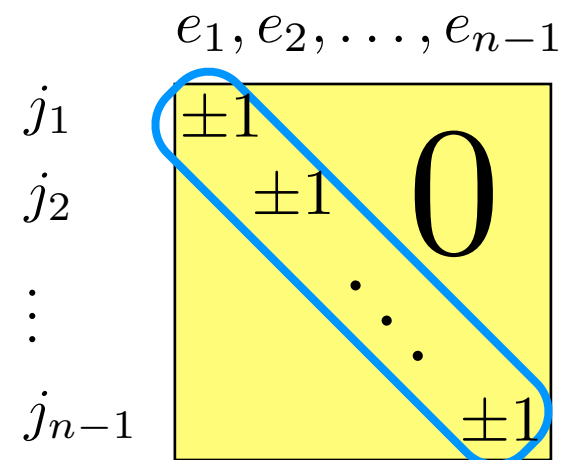
$\exists$  a leaf  $j_1 \neq i$  with incident edge  $e_1$ , delete  $e_1$

$\exists$  a leaf  $j_2 \neq i$  with incident edge  $e_2$ , delete  $e_2$

$\vdots$

vertices:  $j_1, j_2, \dots, j_{n-1}$

edges:  $e_1, e_2, \dots, e_{n-1}$



$\det(B') = \pm 1$

## Cauchy-Binet

$$\det(L_{i,i}) = \sum_{S \in \binom{[m]}{n-1}} \det(B_{[n] \setminus \{i\}, S})^2$$

$$j \in [n] \setminus \{i\}, e \in S$$

$$B_{[n] \setminus \{i\}, S}(j, e) = \begin{cases} 1 & e = \{j, k\}, j < k \\ -1 & e = \{j, k\}, j > k \\ 0 & \text{otherwise} \end{cases}$$

$$\det(B_{[n] \setminus \{i\}, S}) = \begin{cases} \pm 1 & S \text{ is a spanning tree of } G \\ 0 & \text{otherwise} \end{cases}$$

## Kirchhoff's Matrix-Tree Theorem:

$$\forall i, \quad t(G) = \det(L_{i,i})$$

all  $n$ -vertex trees: spanning trees of  $K_n$

$$L_{i,i} = \begin{bmatrix} n-1 & -1 & \cdots & -1 \\ -1 & n-1 & \cdots & -1 \\ \vdots & \vdots & \ddots & \vdots \\ -1 & -1 & \cdots & n-1 \end{bmatrix}$$

Cayley formula:

$$T_n = t(K_n) = \det(L_{i,i}) = n^{n-2}$$