

# Combinatorics

## Extremal Graph Theory

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# Extremal Combinatorics

“how large or how small a collection of finite objects can be, if it has to satisfy certain restrictions”

## Extremal Problem:

“What is the largest number of edges that an  $n$ -vertex *cycle-free* graph can have?”

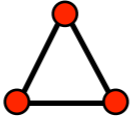
$$(n - 1)$$

## Extremal Graph:

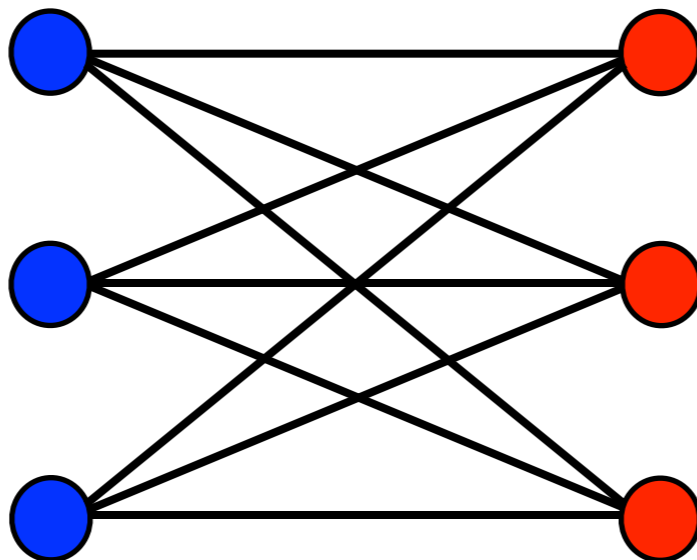
spanning tree

# Triangle-Freeness

# Triangle-free graph

contains no  as subgraph

**Example:** bipartite graph



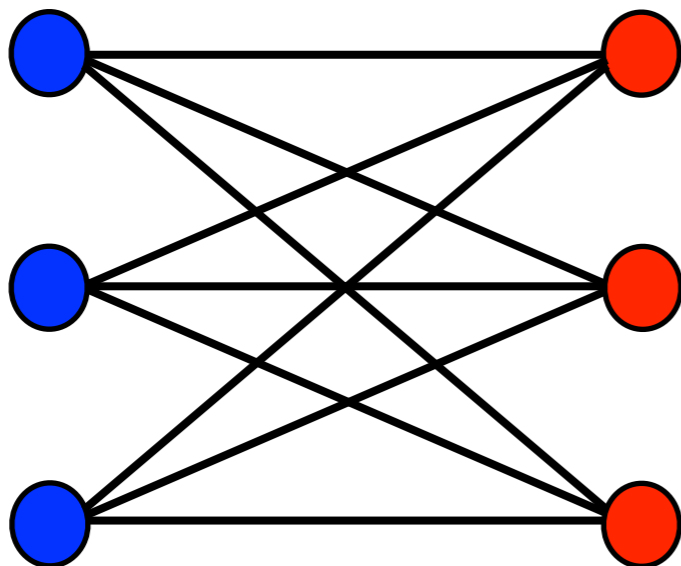
$|E|$  is maximized for  
complete balanced bipartite graph

**Extremal?**

# Mantel's Theorem

## Theorem (Mantel 1907)

If  $G(V, E)$  has  $|V| = n$  and is **triangle-free**,  
then  $|E| \leq \frac{n^2}{4}$ .



For  $n$  is even,  
extremal graph:

$$K_{\frac{n}{2}, \frac{n}{2}}$$

$$\triangle\text{-free} \implies |E| \leq n^2/4$$

**First Proof.** Induction on  $n$ .

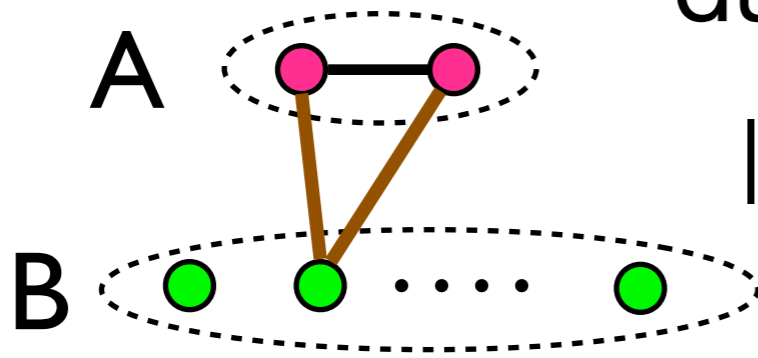
**Basis:**  $n = 1, 2$ . trivial

**Induction Hypothesis:** for any  $n < N$

$$|E| > \frac{n^2}{4} \implies G \supseteq \triangle$$

**Induction step:** for  $n = N$

due to **I.H.**  $|E(B)| \leq (n - 2)^2/4$



$$|E(A, B)| = |E| - |E(B)| - 1$$

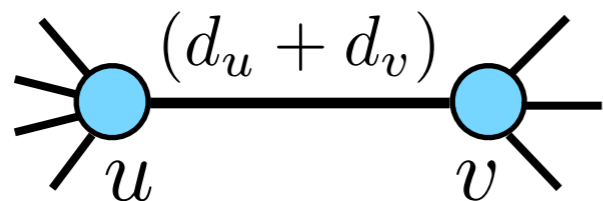
$$> \frac{n^2}{4} - \frac{(n - 2)^2}{4} - 1 = n - 2$$

**pigeonhole!**

$$\triangle\text{-free} \implies |E| \leq n^2/4$$

**Second Proof.**

$\triangle$ -free



$$\implies d_u + d_v \leq n, \quad \forall uv \in E$$



Double counting: 
$$\sum_{v \in V} d_v^2 = \sum_{uv \in E} (d_u + d_v) \leq n |E|$$

**Cauchy-Schwarz**

*(handshaking)*

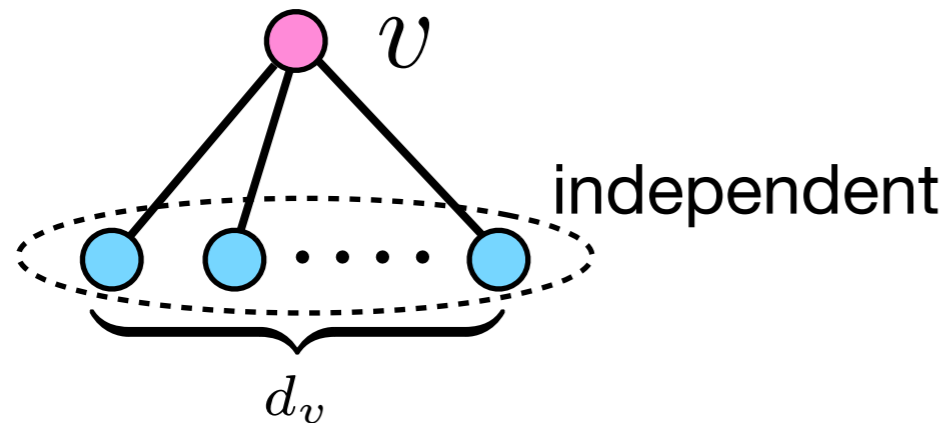
$$n^2 |E| \geq n \sum_{v \in V} d_v^2 = \left( \sum_{v \in V} 1^2 \right) \left( \sum_{v \in V} d_v^2 \right) \geq \left( \sum_{v \in V} d_v \right)^2 = 4 |E|^2$$

$$\implies |E| \leq n^2/4$$

$$\triangle\text{-free} \implies |E| \leq n^2/4$$

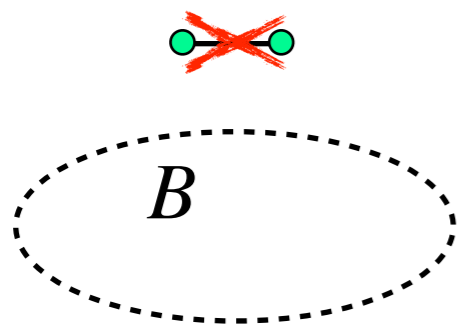
### Third Proof.

**A:** maximum independent set  $\alpha = |A|$



$$\forall v \in V, d_v \leq \alpha$$

$B = V \setminus A$   $B$  incident to all edges  $\beta = |B|$



Inequality of the arithmetic and geometric mean

$$|E| \leq \sum_{v \in B} d_v \leq \alpha \beta \leq \left( \frac{\alpha + \beta}{2} \right)^2 = \frac{n^2}{4}$$

# Turán's Theorem



Paul Turán  
(1910-1976)

# Turán's Theorem

“Suppose  $G$  is a  $K_r$ -free graph.  
What is the largest number of  
edges that  $G$  can have?”

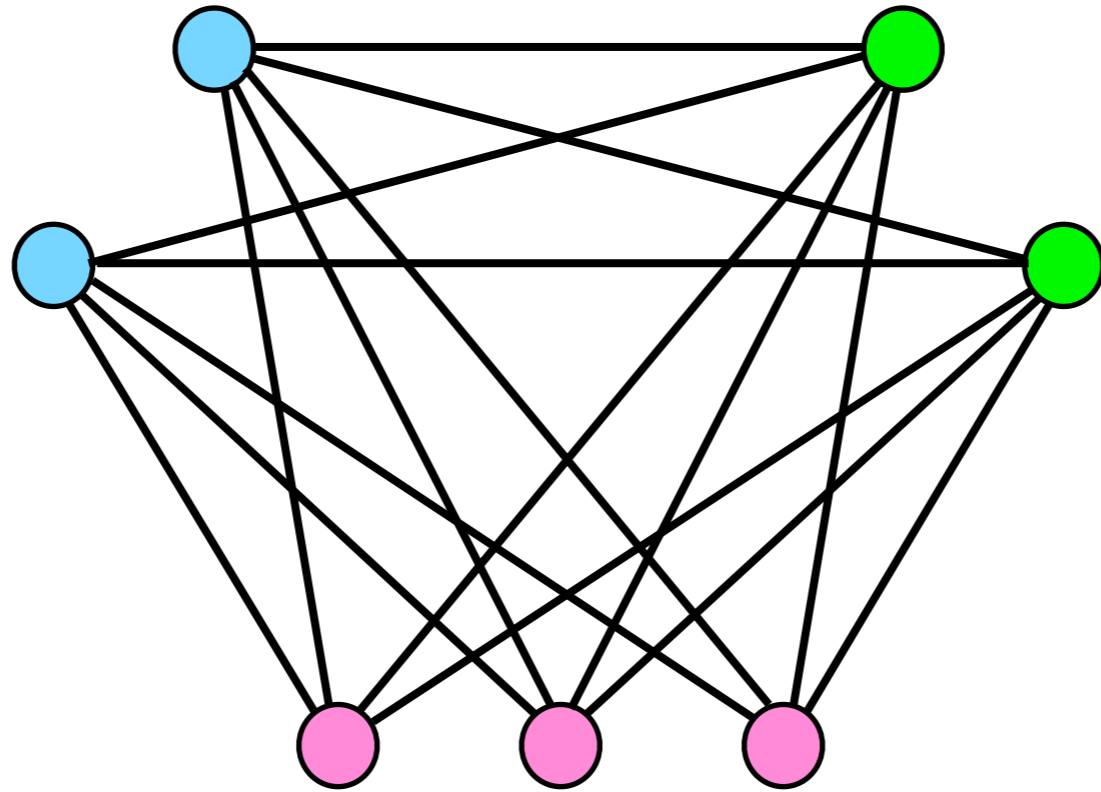
## Theorem (Turán 1941)

If  $G(V, E)$  has  $|V| = n$  and is  $K_r$ -free, then

$$|E| \leq \frac{r-2}{2(r-1)} n^2$$

Complete multipartite graph  $K_{n_1, n_2, \dots, n_r}$

$K_{2,2,3}$



Turán graph  $T(n, r)$ :

$$T(n, r) = K_{n_1, n_2, \dots, n_r}$$

where  $n_1 + \dots + n_r = n$  and  $n_i \in \left\{ \left\lfloor \frac{n}{r} \right\rfloor, \left\lceil \frac{n}{r} \right\rceil \right\}$

Turán graph  $T(n, r)$ :

$$T(n, r) = K_{n_1, n_2, \dots, n_r}$$

where  $n_1 + \dots + n_r = n$  and  $n_i \in \left\{ \left\lfloor \frac{n}{r} \right\rfloor, \left\lceil \frac{n}{r} \right\rceil \right\}$

$T(n, r - 1)$  has no  $K_r$

$$\begin{aligned} |T(n, r - 1)| &\leq \binom{r - 1}{2} \left( \frac{n}{r - 1} \right)^2 \\ &= \frac{r - 2}{2(r - 1)} n^2 \end{aligned}$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)}n^2$$

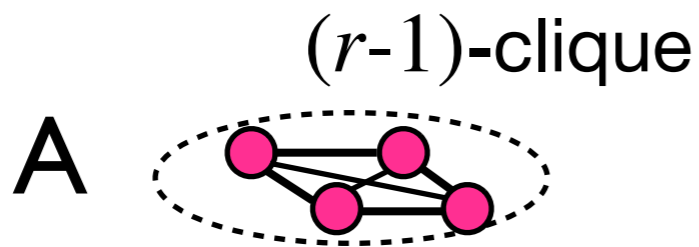
## First Proof. (Induction)

**Basis:**  $n = 1, 2, \dots, r - 1$ .

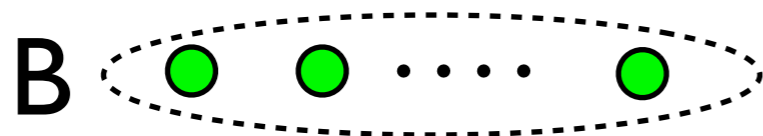
**Induction Hypothesis:** true for any  $n < N$

**Induction step:** for  $n = N$ ,

suppose  $G$  is **maximum  $K_r$ -free**



$\exists (r - 1)$ -clique

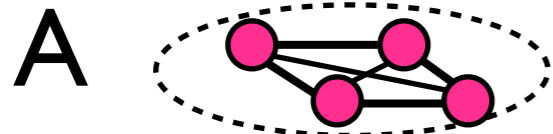


$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)}n^2$$

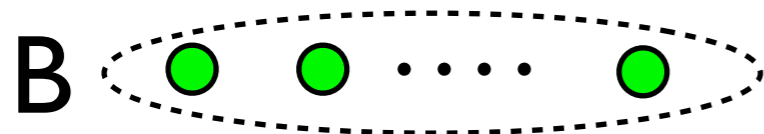
## First Proof. (Induction)

suppose  $G$  is maximum  $K_r$ -free

( $r-1$ )-clique **I.H.:**  $|E(B)| \leq \frac{r-2}{2(r-1)}(n-r+1)^2$



$K_r$ -free  $\implies$  no  $u \in B \sim$  all  $v \in A$



$\implies |E(A, B)| \leq (r-2)(n-r+1)$

$$|E| = |E(A)| + |E(B)| + |E(A, B)|$$

$$\leq \binom{r-1}{2} + \frac{r-2}{2(r-1)}(n-r+1)^2 + (r-2)(n-r+1)$$

$$= \frac{r-2}{2(r-1)}n^2$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)} n^2$$

## Second Proof. (weight shifting)

Assign each vertex  $v$  a **weight**  $w_v > 0$  s.t.  $\sum_{v \in V} w_v = 1$

$$\text{Evaluate } S(\vec{w}) = \sum_{uv \in E} w_u w_v$$

$$\text{Let } W_u = \sum_{v \sim u} w_v \quad \text{For } u \approx v \text{ that } W_u \geq W_v$$

$$(w_u + \epsilon)W_u + (w_v - \epsilon)W_v \geq w_u W_u + w_v W_v$$

**shifting** all weight of  $v$  to  $u \implies S(\vec{w})$  non-decreasing

$S(\vec{w})$  is maximized  $\implies$  all weights on a clique

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)}n^2$$

## Second Proof. (weight shifting)

Assign each vertex  $v$  a **weight**  $w_v > 0$  s.t.  $\sum_{v \in V} w_v = 1$

$$\text{Evaluate } S(\vec{w}) = \sum_{uv \in E} w_u w_v \leq \binom{r-1}{2} \frac{1}{(r-1)^2}$$

$S(\vec{w})$  is maximized  $\implies$  all weights on a clique

$$\text{when all } w_v = \frac{1}{n}$$

$$S(\vec{w}) = \sum_{uv \in E} w_u w_v = \frac{|E|}{n^2}$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)} n^2$$

**Third Proof.** (The probabilistic method)

clique number  $\omega(G)$ : size of the largest clique

$$\omega(G) \geq \sum_{v \in V} \frac{1}{n - d_v}$$

random permutation  $\pi$  of  $V$

$S = \{v \mid \pi_u < \pi_v \implies u \sim v\}$   
is a clique

**Linearity of expectation:**

$$\begin{aligned} \mathbb{E}[|S|] &= \sum_{v \in V} \Pr[v \in S] \geq \sum_{v \in V} \Pr[\forall u \not\sim v : \pi_u \geq \pi_v] \\ &= \sum_{v \in V} \frac{1}{n - d_v} \end{aligned}$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)}n^2$$

**Third Proof.** (The probabilistic method)

$$\omega(G) \geq \sum_{v \in V} \frac{1}{n - d_v}$$

**Cauchy-Schwarz**

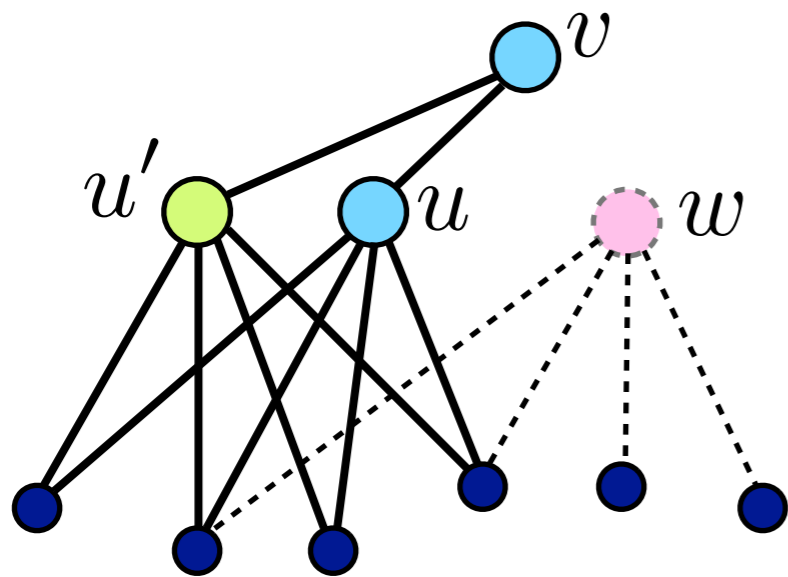
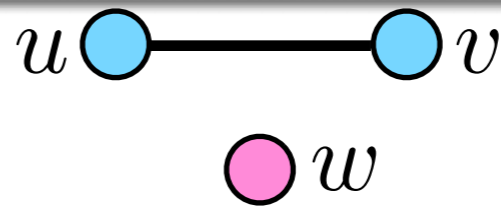
$$\begin{aligned} n^2 &= \left( \sum_{v \in V} 1 \right)^2 \leq \left( \sum_{v \in V} \frac{1}{n - d_v} \right) \left( \sum_{v \in V} (n - d_v) \right) \\ &\leq \omega(G) \sum_{v \in V} (n - d_v) = (r-1)(n^2 - 2|E|) \\ &\quad \text{(handshaking)} \\ &\implies |E| \leq \frac{r-2}{2(r-1)}n^2 \end{aligned}$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)} n^2$$

## Fourth Proof.

Suppose  $G$  is  $K_r$ -free with **maximum** edges.

$G$  does not have



By contradiction.

**Case.1**  $d_w < d_u$  or  $d_w < d_v$

duplicate  $u$ , delete  $w$ , **still  $K_r$ -free**

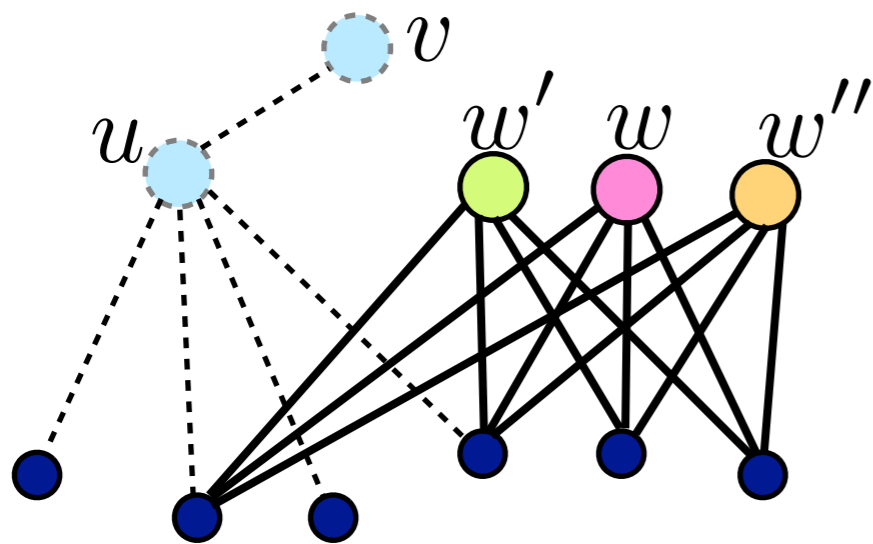
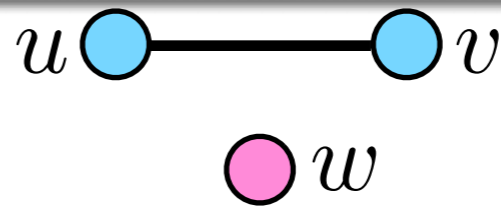
$$|E'| = |E| + d_u - d_w > |E|$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)} n^2$$

## Fourth Proof.

Suppose  $G$  is  $K_r$ -free with **maximum** edges.

$G$  does not have



**Case.2**  $d_w \geq d_u \wedge d_w \geq d_v$

delete  $u, v$ , duplicate  $w$ , twice

still  $K_r$ -free

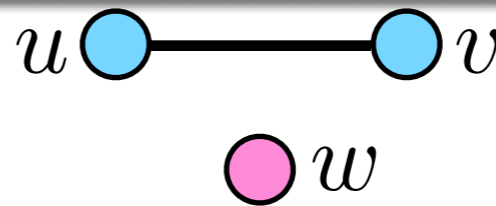
$$|E'| = |E| + 2d_w - (d_u + d_v - 1) > |E|$$

$$K_r\text{-free} \implies |E| \leq \frac{r-2}{2(r-1)}n^2$$

## Fourth Proof.

Suppose  $G$  is  $K_r$ -free with **maximum** edges.

$G$  does not have



$u \approx v$  is an equivalence relation

$G$  is a complete multipartite graph

optimize  $K_{n_1, n_2, \dots, n_{r-1}}$

subject to  $n_1 + n_2 + \dots + n_{r-1} = n$

## Turán's Theorem (clique)

If  $G(V, E)$  has  $|V| = n$  and is  $K_r$ -free, then

$$|E| \leq \frac{r-2}{2(r-1)} n^2$$

## Turán's Theorem (independent set)

If  $G(V, E)$  has  $|V| = n$  and  $|E| = m$ , then  $G$  has an independent set of size

$$\geq \frac{n^2}{2m + n}$$

# Parallel Max

- compute max of  $n$  distinct numbers
- computation model: **parallel, comparison-based**
- 1-round algorithm:  $\binom{n}{2}$  comparisons of all pairs
- lower bound for one-round:
  - $\binom{n}{2}$  comparisons are required in the worst case



**adversary argument**

# Parallel Max

- 2-round algorithm:

- divide  $n$  numbers into  $k$  groups of  $n/k$  each

- *1st round*: find max of each group;

$$k \binom{n/k}{2} \text{ comparisons}$$

- *2nd round*: find the max of the  $k$  maxes

$$\binom{k}{2} \text{ comparisons}$$

- **total comparisons:**  $k \binom{n/k}{2} + \binom{k}{2} = O\left(n^{4/3}\right)$

for  $k = n^{2/3}$

3-round?

optimal?

1st round:

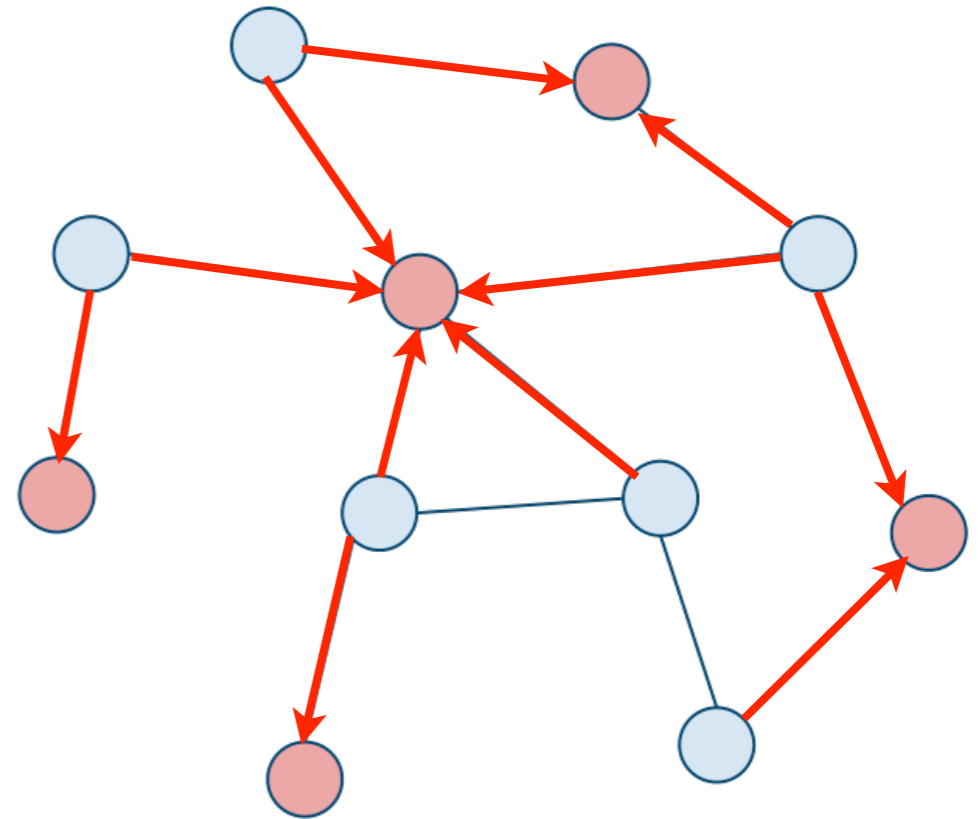
Alg:  $m$  comparisons



choose an independent set

of size  $\geq \frac{n^2}{2m+n}$  (Turán)

make them local maximal



2nd round:

a parallel max problem of size  $\geq \frac{n^2}{2m+n}$

requires  $\geq \binom{\frac{n^2}{2m+n}}{2}$  comparisons

total comparisons  $\geq m + \binom{\frac{n^2}{2m+n}}{2} = \Omega(n^{4/3})$

# ***Fundamental Theorem*** **of Extremal Graph Theory**

# Extremal Graph Theory

Fix a graph  $H$ .

$$\text{ex}(n, H)$$

largest possible number of edges  
of  $G \not\supseteq H$  on  $n$  vertices

$$\text{ex}(n, H) = \max_{\substack{G \not\supseteq H \\ |V(G)|=n}} |E(G)|$$

## Turán's Theorem

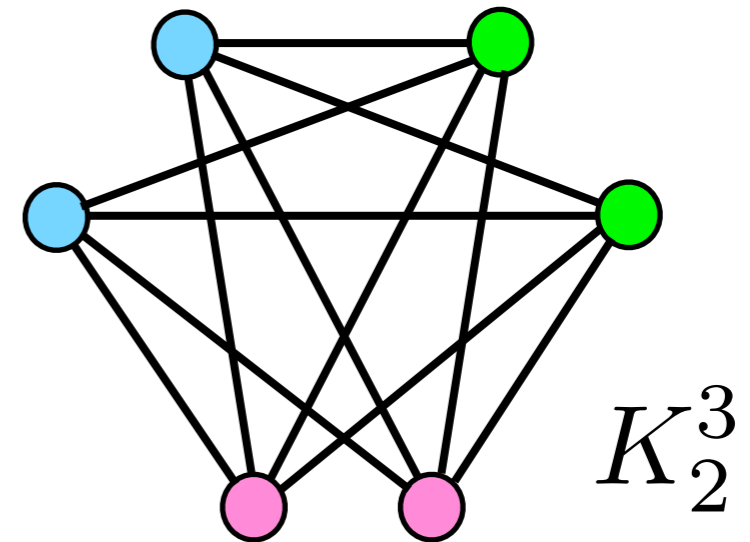
$$\text{ex}(n, K_r) = |T(n, r-1)| \leq \frac{r-2}{2(r-1)} n^2$$

# Erdős–Stone theorem

(Fundamental theorem of extremal graph theory)

$$K_s^r = K_{\underbrace{s, s, \dots, s}_r} = T(rs, r)$$

complete  $r$ -partite graph  
with  $s$  vertices in each part



**Theorem (Erdős–Stone 1946)**

$$\text{ex}(n, K_s^r) = \left( \frac{r-2}{2(r-1)} + o(1) \right) n^2$$

## Theorem (Erdős–Stone 1946)

$$\text{ex}(n, K_s^r) = \left( \frac{r-2}{2(r-1)} + o(1) \right) n^2$$

$\text{ex}(n, H) / \binom{n}{2}$  **extremal density** of subgraph  $H$

## Corollary

For any nonempty graph  $H$

$$\lim_{n \rightarrow \infty} \frac{\text{ex}(n, H)}{\binom{n}{2}} = \frac{\chi(H) - 2}{\chi(H) - 1}$$

$$\lim_{n \rightarrow \infty} \frac{\text{ex}(n, H)}{\binom{n}{2}} = \frac{\chi(H) - 2}{\chi(H) - 1}$$

$$\chi(H) = r$$

$H \not\subseteq T(n, r - 1)$  for any  $n$

$$\text{ex}(n, H) \geq |T(n, r - 1)|$$

$H \subseteq K_s^r$  for sufficiently large  $s$

$$\text{ex}(n, H) \leq \text{ex}(n, K_s^r)$$

$$= \left( \frac{r-2}{2(r-1)} + o(1) \right) n^2$$

$$\lim_{n \rightarrow \infty} \frac{\text{ex}(n, H)}{\binom{n}{2}} = \frac{\chi(H) - 2}{\chi(H) - 1}$$

$$\chi(H) = r$$

$$|T(n, r - 1)| \leq \text{ex}(n, H) \leq \left( \frac{r - 2}{2(r - 1)} + o(1) \right) n^2$$

$$\frac{r - 2}{r - 1} - o(1) \leq \frac{\text{ex}(n, H)}{\binom{n}{2}} \leq \frac{r - 2}{r - 1} + o(1)$$

# Cycles


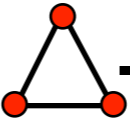
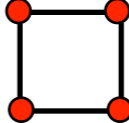
# Girth

**girth  $g(G)$ :** length of the shortest cycle in  $G$

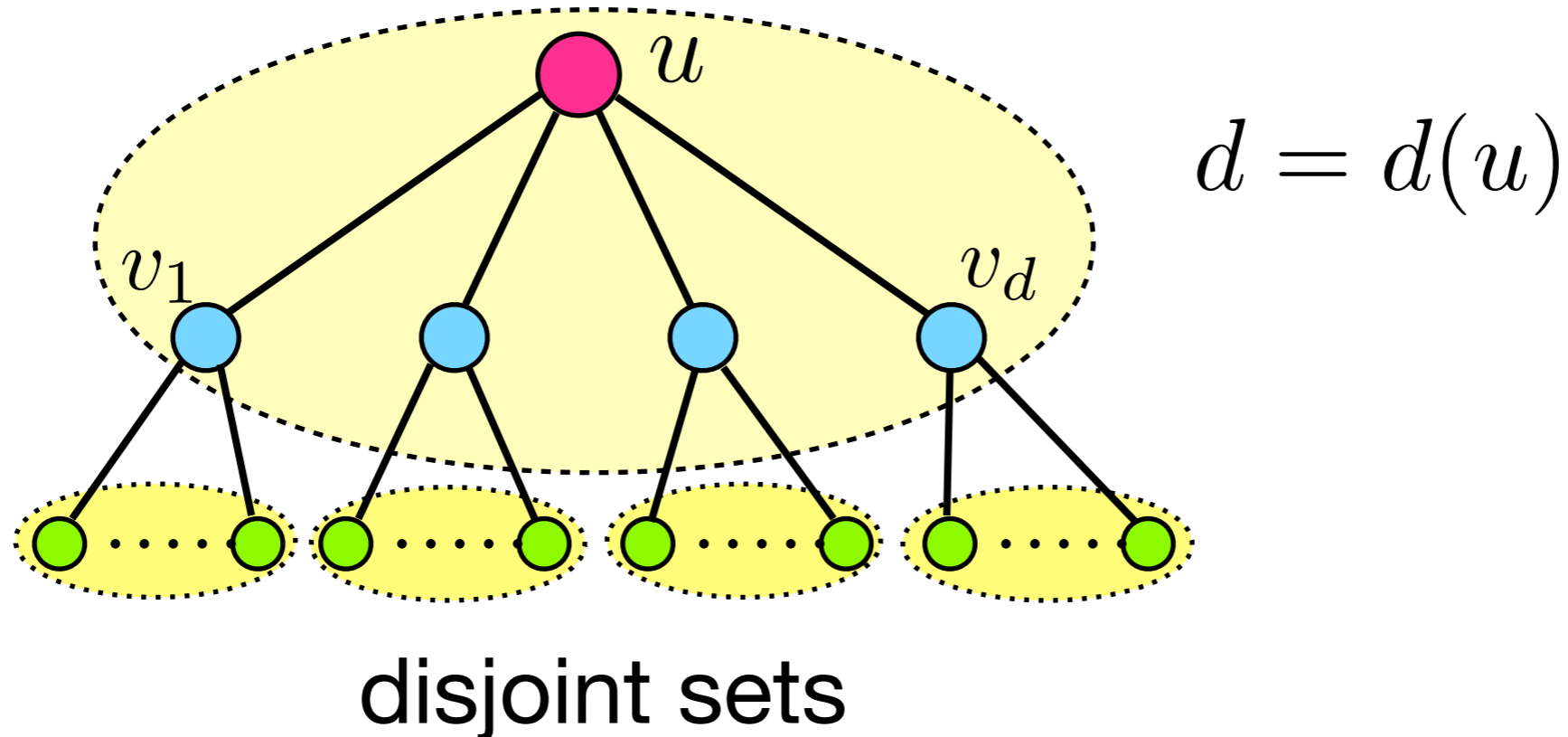
## Theorem

If  $G(V, E)$  has  $|V| = n$  and  $\text{girth } g(G) \geq 5$ ,

$$\text{then } |E| \leq \frac{1}{2}n\sqrt{n-1}$$

$g(G) \geq 5$   - and -free

$$g(G) \geq 5 \Rightarrow |E| \leq \frac{1}{2}n\sqrt{n-1}$$

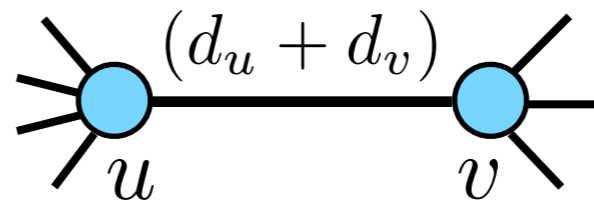


$$(d + 1) + (d(v_1) - 1) + \cdots + (d(v_d) - 1) \leq n$$

$$\sum_{v:v \sim u} d(v) \leq n - 1$$

$$g(G) \geq 5 \Rightarrow |E| \leq \frac{1}{2}n\sqrt{n-1}$$

$$\forall u \in V, \sum_{v:v \sim u} d(v) \leq n-1$$



$$\begin{aligned} n(n-1) &\geq \sum_{u \in V} \sum_{v:v \sim u} d(v) = \sum_{v \in V} d(v)^2 \\ &\geq \frac{\left(\sum_{v \in V} d(v)\right)^2}{n} = \frac{4|E|^2}{n} \end{aligned}$$

**Cauchy-Schwarz**

# Hamiltonian Cycle

## Dirac's Theorem

$\forall v \in V, d_v \geq \frac{n}{2} \Rightarrow G(V, E)$  is Hamiltonian.

By contradiction, suppose  $G$  is the **maximum non-Hamiltonian** graph with  $\forall v \in V, d_v \geq \frac{n}{2}$

**adding 1 edge  $\implies$  Hamiltonian**

$\exists$  a Hamiltonian path

say  $v_1 v_2 \cdots v_n$

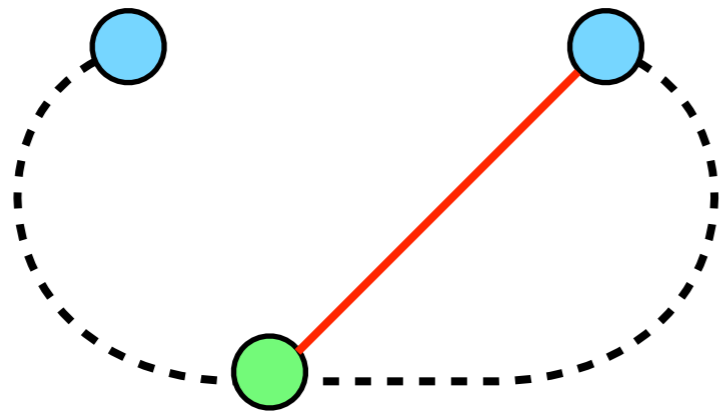
$G$  is non-Hamiltonian

$$\forall v \in V, d_v \geq \frac{n}{2}$$

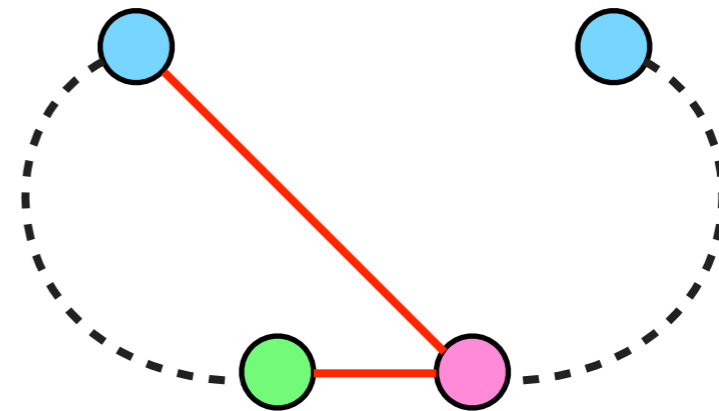
$\exists$  a Hamiltonian path

$$v_1 v_2 \cdots v_n$$

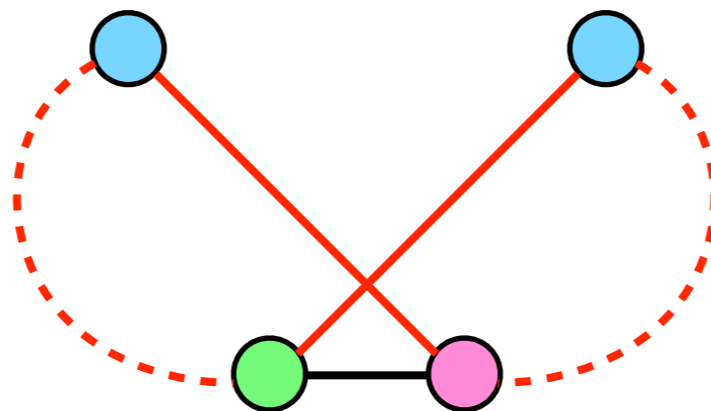
$$\{i \mid v_i \sim v_n\}$$



$$\{i \mid v_{i+1} \sim v_1\}$$



$\geq \frac{n}{2} + \frac{n}{2}$  pigeons in  $\{1, 2, \dots, n-1\}$



**Contradiction!**