## Homework 7

Course: Algorithm Design and Analysis
Semester: Spring 2024
Instructor: Shi Li
Due Date: 2024/6/12

Student Name: $\qquad$ Student ID: $\qquad$

| Problems | 1 | 2 | 3 | 4 | 5 | Total |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Max. Score | 20 | 16 | 14 | 20 | 30 | 100 |
| Your Score |  |  |  |  |  |  |

Problem 1. For each of the following problems, state (i) whether the problem is provably in NP, and (ii) whether the problem is provably in Co-NP. If you claim a problem is in NP, then you need to describe the certificate and the certifier for the proof.
(1a) Given a graph $G=(V, E)$ with edge weights $w \in \mathbb{Z}_{\geq 0}^{E}$ and an integer $W \geq 0$, the problems asks if there is a spanning tree of $G$ with total weight at most $W$.
(1b) Given two boolean formulas, the problem asks whether the two boolean formulas are equivalent. For example, $\left(x_{1} \vee x_{2}\right) \wedge\left(\neg x_{1} \vee x_{3}\right)$ and $\left(\neg x_{1} \wedge x_{2}\right) \vee\left(x_{1} \wedge x_{3}\right)$ are equivalent since they give the same value for every assignment of $\left(x_{1}, x_{2}, x_{3}\right)$.
(1c) An undirected graph $G=(V, E)$ is called an expander if for every $S \subseteq V$, the number of edges between $S$ and $V \backslash S$ in $G$ is at least $\min \{|S|,|V \backslash S|\}$. Given a graph $G$, the problem asks if $G$ is an expander or not.
(1d) Given $n$ items [ $n$ ] with integer weights $w_{1}, w_{2}, \cdots, w_{n} \geq 0$ and integer values $v_{1}, v_{2}, \cdots, v_{n} \geq 0$, and two integers $W$ and $V$, the problem asks if there is a set $S \subseteq[n]$ of items with total weight at most $W$ and total value at least $V$.

Problem 2. Indicate if each of the following statements is true or false. A true/false answer for each statement is sufficient; you do not need to give proofs/counterexamples for your answers.
(2a) If a decision problem $X$ can be solved in polynomial time, then $X \notin$ NP.
(2b) Assume $X$ is a NP-complete problem and $X$ has a polynomial time algorithm. Then $P=N P$.
(2c) If a problem $X$ is NP-compelte, then the circuit-satisfiability problem is polynomialtime reducible to $X$.
(2d) Based on our knowledge, it is possible that $\mathrm{P} \cap \mathrm{NP}=\emptyset$.

Problem 3. Prove that $\mathrm{P}=\mathrm{NP}$ if and only if $\mathrm{P}=\mathrm{Co}-\mathrm{NP}$.
Problem 4. In class, we proved that HP (Hamiltonian Path) $\leq_{P}$ HC (Hamiltonian Cycle). Prove the other direction, i.e, $\mathrm{HC} \leq_{P} \mathrm{HP}$.

Problem 5. In the Steiner Tree problem, we are given a graph $G=(V, E)$, with edge weights $w \in \mathbb{Z}_{\geq 0}^{E}$, and a set $X \subseteq V$ of vertices. The goal of the problem is to find the minimum-weight edges to connect $X$ in the graph $G$. Formally, our goal is to find a tree $T=\left(V_{T}, E_{T}\right)$ such that $X \subseteq V_{T} \subseteq V$ and $E_{T} \subseteq E$ (such a tree $T$ is called a Steiner tree for $X$ ), so as to minimize $\sum_{e \in E_{t}} w_{e}$.

In the decision problem, we are additionally given an integer bound $W>0$, and we need to decide if there is a Steiner tree for $X$ with weight at most $W$.

Prove that Vertex-Cover $\leq_{P}$ Steiner-Tree.

