

算法设计与分析(2024年春季学期)

Dynamic Programming

授课老师: 栗师

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Paradigms for Designing Algorithms

Greedy algorithm

- Make a greedy choice
- Prove that the greedy choice is safe
- Reduce the problem to a sub-problem and solve it iteratively
- Usually for optimization problems

Divide-and-conquer

- Break a problem into many **independent** sub-problems
- Solve each sub-problem separately
- Combine solutions for sub-problems to form a solution for the original one
- Usually used to design more efficient algorithms

Paradigms for Designing Algorithms

Dynamic Programming

- Break up a problem into many **overlapping** sub-problems
- Build solutions for larger and larger sub-problems
- Use a **table** to store solutions for sub-problems for reuse

Recall: Computing the n -th Fibonacci Number

- $F_0 = 0, F_1 = 1$
- $F_n = F_{n-1} + F_{n-2}, \forall n \geq 2$
- Fibonacci sequence: 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, \dots

Fib(n)

```
1:  $F[0] \leftarrow 0$   
2:  $F[1] \leftarrow 1$   
3: for  $i \leftarrow 2$  to  $n$  do  
4:    $F[i] \leftarrow F[i - 1] + F[i - 2]$   
5: return  $F[n]$ 
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- Store each $F[i]$ for future use.

Outline

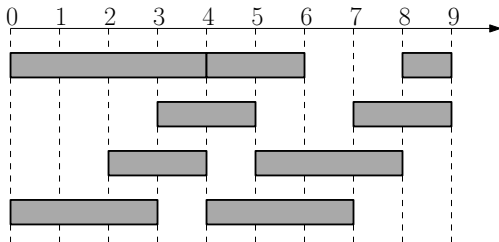
- 1 Weighted Interval Scheduling
- 2 Segmented Least Squares
- 3 Subset Sum Problem
 - Related Problem: Knapsack Problem
- 4 Longest Common Subsequence
 - Longest Common Subsequence in Linear Space
- 5 Shortest Paths in Directed Acyclic Graphs
- 6 Matrix Chain Multiplication
- 7 Optimum Binary Search Tree
- 8 Summary

Recall: Interval Scheduling

Input: n jobs, job i with start time s_i and finish time f_i

i and j are compatible if $[s_i, f_i)$ and $[s_j, f_j)$ are disjoint

Output: a maximum-size subset of mutually compatible jobs

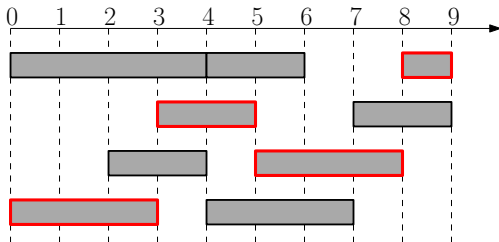


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Weighted Interval Scheduling

Input: n jobs, job i with start time s_i and finish time f_i

each job has a weight (or value) $v_i > 0$

i and j are compatible if $[s_i, f_i)$ and $[s_j, f_j)$ are disjoint

Output: a **maximum-weight** subset of mutually compatible jobs

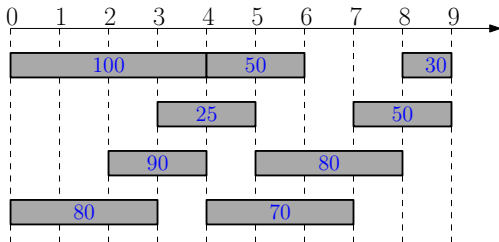
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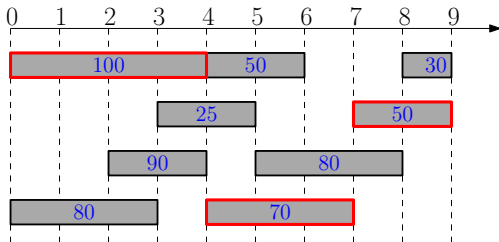
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Optimum value = 220

Hard to Design a Greedy Algorithm

Q: Which job is safe to schedule?

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- Job with the earliest finish time?

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- Job with the earliest finish time? No, we are ignoring weights
- Job with the largest weight?

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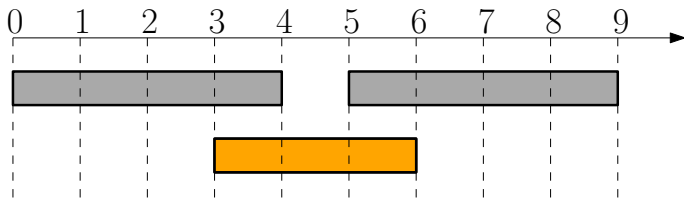
No, when weights are equal, this is the shortest job

Hard to Design a Greedy Algorithm

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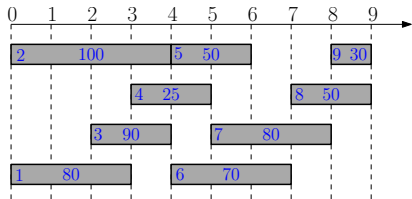
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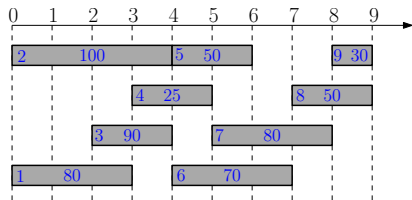
Designing a Dynamic Programming Algorithm

Designing a Dynamic Programming Algorithm



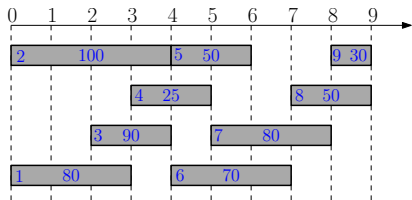
- Sort jobs according to non-decreasing order of finish times

Designing a Dynamic Programming Algorithm



- Sort jobs according to non-decreasing order of finish times
- $opt[i]$: optimal value for instance only containing jobs $\{1, 2, \dots, i\}$

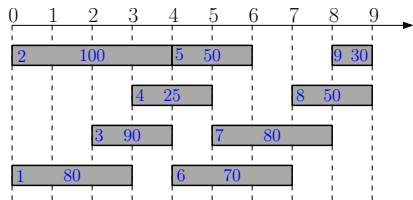
Designing a Dynamic Programming Algorithm



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i	$opt[i]$
0	
1	
2	
3	
4	
5	
6	
7	
8	
9	

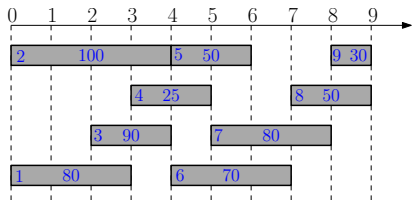
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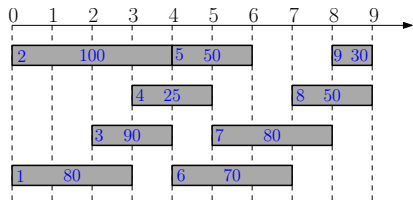
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i	$opt[i]$
0	0
1	80
2	
3	
4	
5	
6	
7	
8	
9	

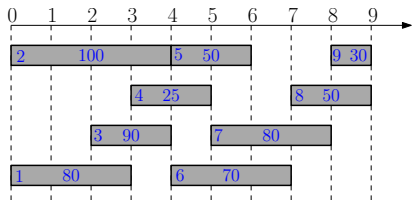
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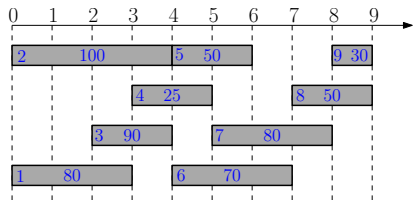
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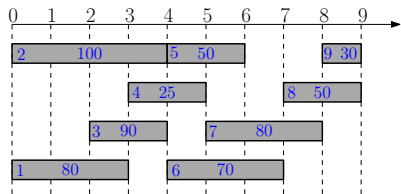
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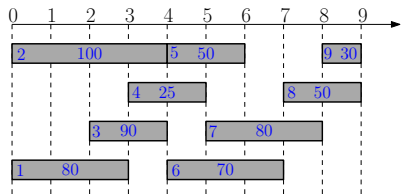
i	$opt[i]$
0	0
1	80
2	100
3	100
4	105
5	150
6	170
7	185
8	220
9	220

Designing a Dynamic Programming Algorithm



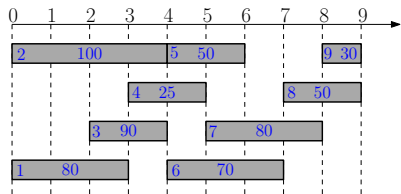
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Designing a Dynamic Programming Algorithm



- Focus on instance $\{1, 2, 3, \dots, i\}$,
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- assume we have computed $opt[0], opt[1], \dots, opt[i - 1]$

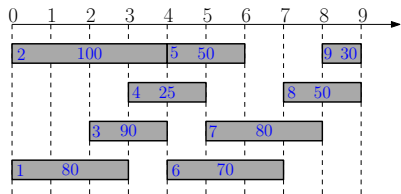
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Q: The value of optimal solution that **does not contain** i ?

Designing a Dynamic Programming Algorithm

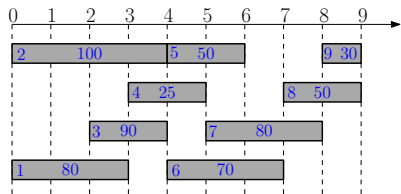


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Designing a Dynamic Programming Algorithm



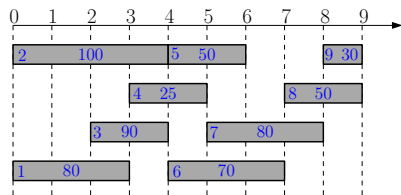
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Q: The value of optimal solution that **contains** job i ?

Designing a Dynamic Programming Algorithm



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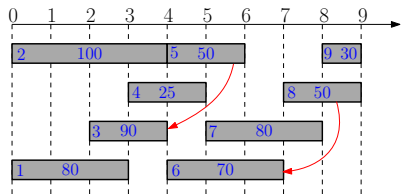
A: $opt[i - 1]$

Q: The value of optimal solution that **contains** job i ?

A: $v_i + opt[p_i]$,

$p_i =$ the largest j such that $f_j \leq s_i$

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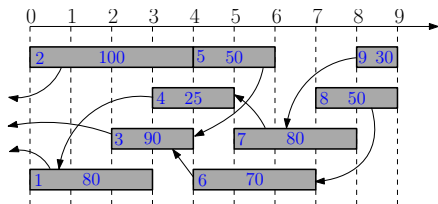
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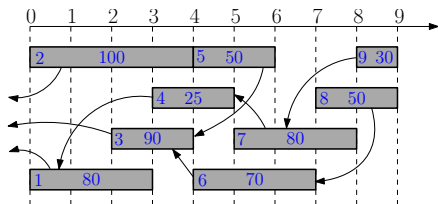


- $opt[0] = 0$
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- $opt[3] =$
- $opt[4] =$
- $opt[5] =$

Designing a Dynamic Programming Algorithm

Recursion for $opt[i]$:

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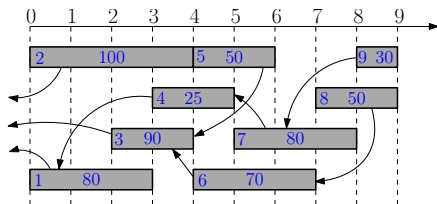


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Designing a Dynamic Programming Algorithm

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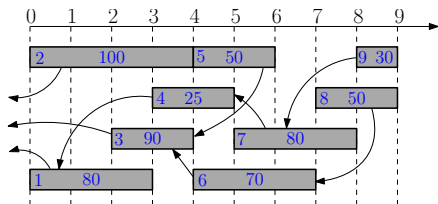


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Designing a Dynamic Programming Algorithm

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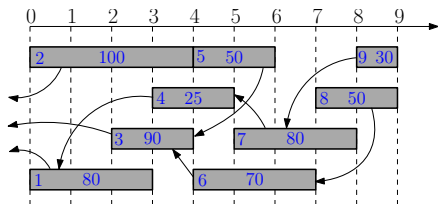


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Designing a Dynamic Programming Algorithm

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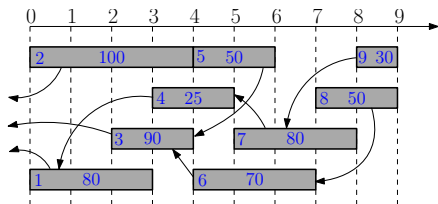


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- $opt[2] = \max\{opt[1], 100 + opt[0]\} = 100$
- $opt[3] = \max\{opt[2], 90 + opt[0]\}$
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- $opt[5] =$

Designing a Dynamic Programming Algorithm

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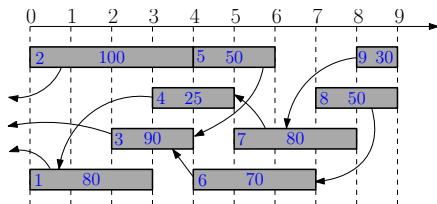


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Designing a Dynamic Programming Algorithm

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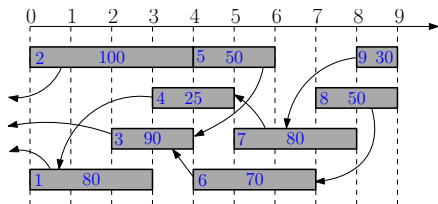


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Designing a Dynamic Programming Algorithm

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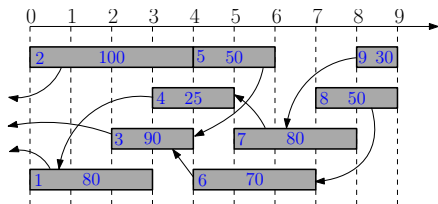


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- $opt[2] = \max\{opt[1], 100 + opt[0]\} = 100$
- $opt[3] = \max\{opt[2], 90 + opt[0]\} = 100$
- $opt[4] = \max\{opt[3], 25 + opt[1]\} = 105$
- $opt[5] =$

Designing a Dynamic Programming Algorithm

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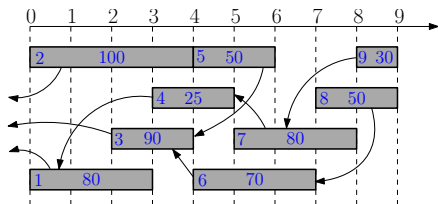


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- $opt[4] = \max\{opt[3], 25 + opt[1]\} = 105$
- $opt[5] = \max\{opt[4], 50 + opt[3]\}$

Designing a Dynamic Programming Algorithm

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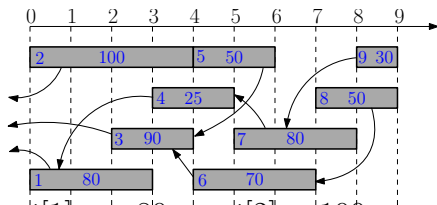


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Designing a Dynamic Programming Algorithm

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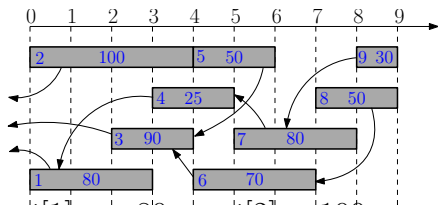


- $opt[0] = 0$, $opt[1] = 80$, $opt[2] = 100$
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Designing a Dynamic Programming Algorithm

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- $opt[0] = 0$, $opt[1] = 80$, $opt[2] = 100$
- $opt[3] = 100$, $opt[4] = 105$, $opt[5] = 150$
- $opt[6] = \max\{opt[5], 70 + opt[3]\} = 170$
- $opt[7] = \max\{opt[6], 80 + opt[4]\} = 185$
- $opt[8] = \max\{opt[7], 50 + opt[6]\} = 220$
- $opt[9] = \max\{opt[8], 30 + opt[7]\} = 220$

Dynamic Programming

- 1: sort jobs by non-decreasing order of finishing times
- 2: compute p_1, p_2, \dots, p_n
- 3: $opt[0] \leftarrow 0$
- 4: **for** $i \leftarrow 1$ to n **do**
- 5: $opt[i] \leftarrow \max\{opt[i - 1], v_i + opt[p_i]\}$

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- Running time sorting: $O(n \lg n)$
- Running time for computing p : $O(n \lg n)$ via binary search
- Running time for computing $opt[n]$: $O(n)$

How Can We Recover the Optimum Schedule?

```
1: sort jobs by non-decreasing order of
   finishing times
2: compute  $p_1, p_2, \dots, p_n$ 
3:  $opt[0] \leftarrow 0$ 
4: for  $i \leftarrow 1$  to  $n$  do
5:     if  $opt[i - 1] \geq v_i + opt[p_i]$  then
6:          $opt[i] \leftarrow opt[i - 1]$ 
7:
8:     else
9:          $opt[i] \leftarrow v_i + opt[p_i]$ 
10:
```

How Can We Recover the Optimum Schedule?

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1: sort jobs by non-decreasing order of
   finishing times
2: compute  $p_1, p_2, \dots, p_n$ 
3:  $opt[0] \leftarrow 0$ 
4: for  $i \leftarrow 1$  to  $n$  do
5:     if  $opt[i - 1] \geq v_i + opt[p_i]$  then
6:          $opt[i] \leftarrow opt[i - 1]$ 
7:          $b[i] \leftarrow N$ 
8:     else
9:          $opt[i] \leftarrow v_i + opt[p_i]$ 
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```

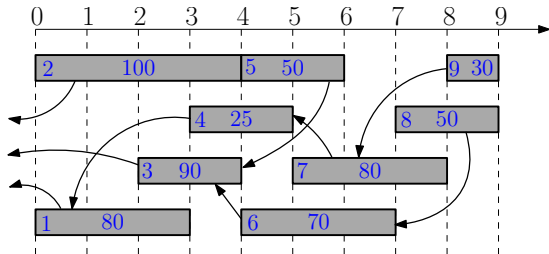
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```

```
1:  $i \leftarrow n, S \leftarrow \emptyset$ 
2: while  $i \neq 0$  do
3:     if  $b[i] = N$  then
4:          $i \leftarrow i - 1$ 
5:     else
6:          $S \leftarrow S \cup \{i\}$ 
7:          $i \leftarrow p_i$ 
8: return  $S$ 
```

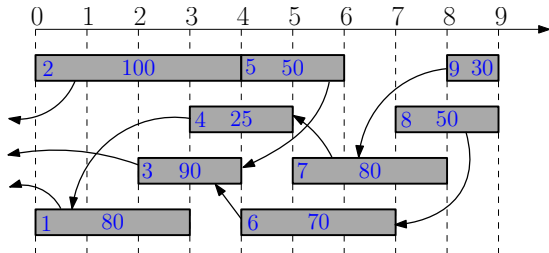
Recovering Optimum Schedule: Example

i	$opt[i]$	$b[i]$
0	0	\perp
1	80	
2	100	
3	100	
4	105	
5	150	
6	170	
7	185	
8	220	
9	220	



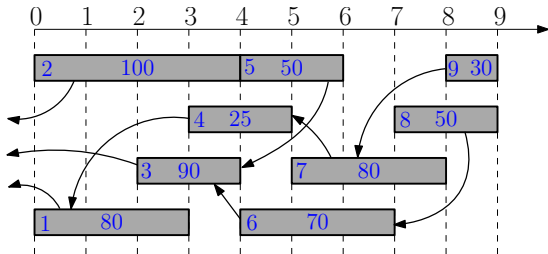
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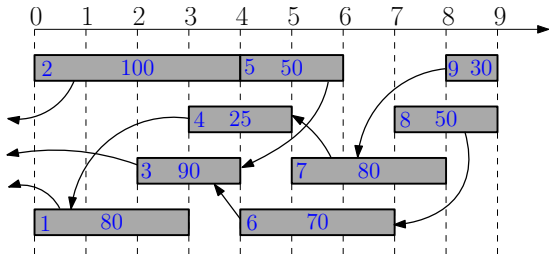
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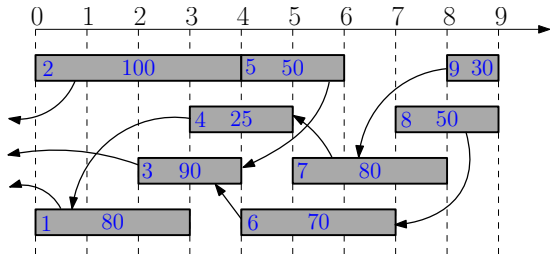
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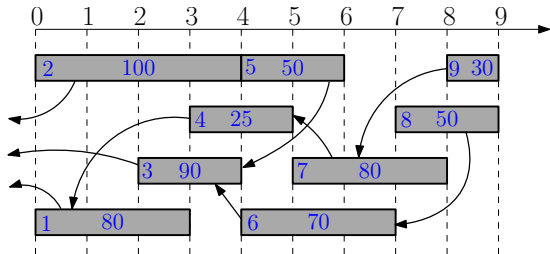
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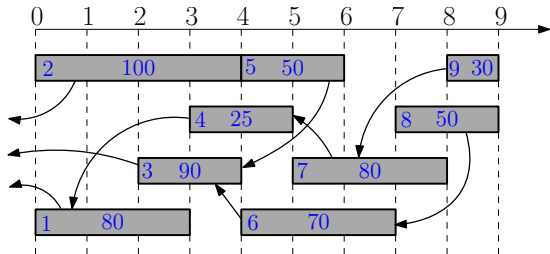
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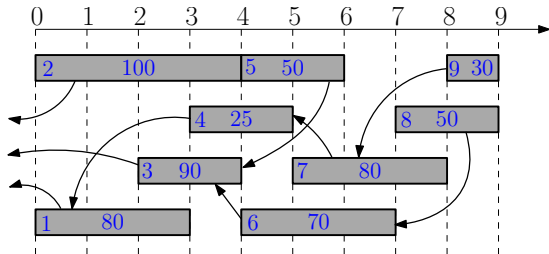
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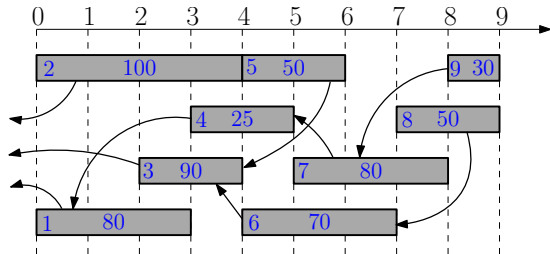
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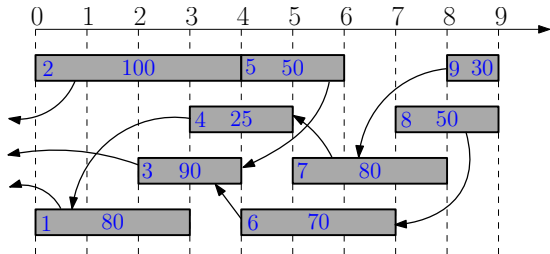
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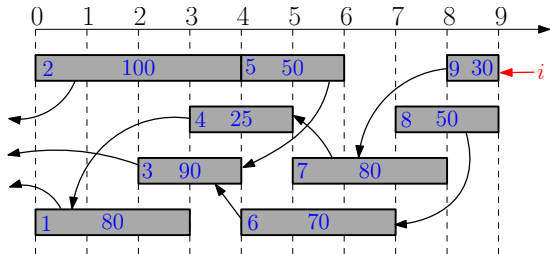
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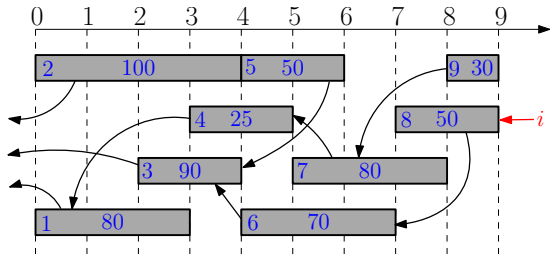
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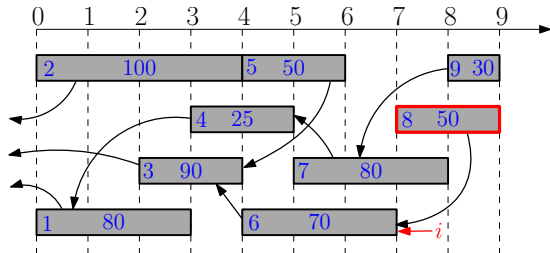
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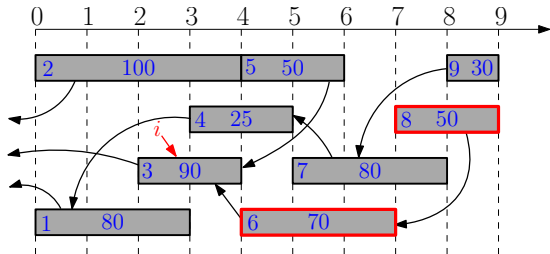
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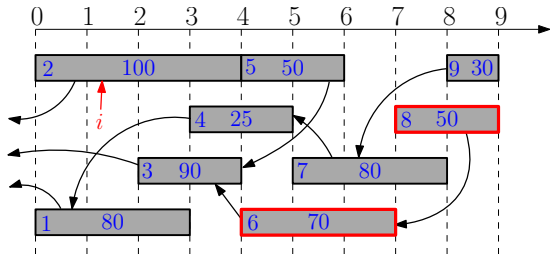
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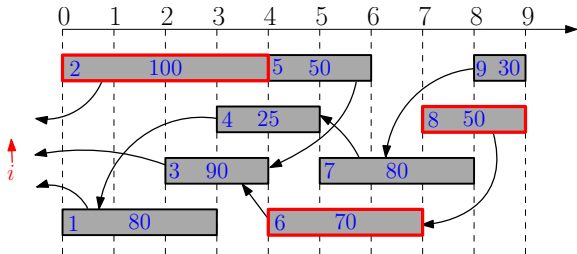
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Dynamic Programming

- Break up a problem into many **overlapping** sub-problems
- Build solutions for larger and larger sub-problems
- Use a **table** to store solutions for sub-problems for reuse

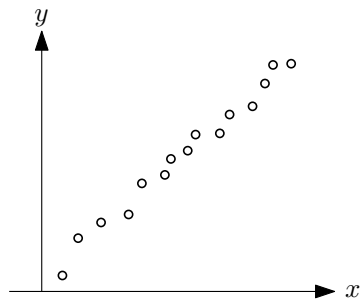
Outline

- 1 Weighted Interval Scheduling
- 2 Segmented Least Squares**
- 3 Subset Sum Problem
 - Related Problem: Knapsack Problem
- 4 Longest Common Subsequence
 - Longest Common Subsequence in Linear Space
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Linear Regression

- $P = \{(x_1, y_1), (x_2, y_2), \dots, (x_n, y_n)\}, x_1 < x_2 < \dots < x_n$
- $L : y = ax + b$

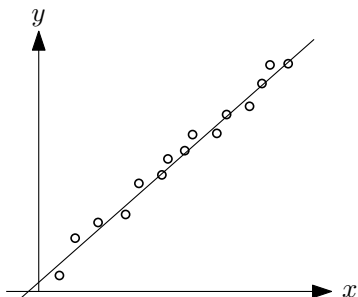
$$\text{Error}(L, P) = \sum_{i=1}^n (y_i - ax_i - b)^2$$



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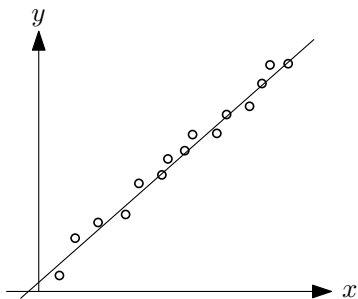
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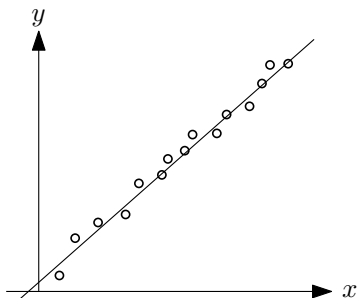
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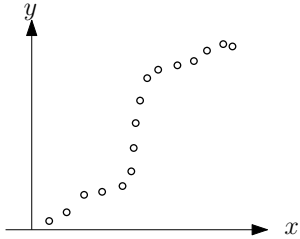


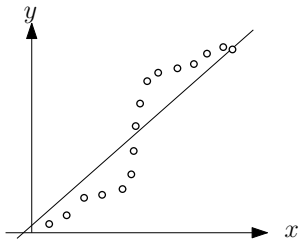
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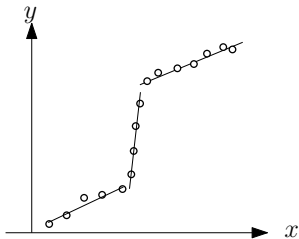
$$a := \frac{n \sum_i x_i y_i - (\sum_i x_i)(\sum_i y_i)}{n \sum_i x_i^2 - (\sum_i x_i)^2}$$

$$b := \frac{\sum_i y_i - a \sum_i x_i}{n}$$





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- Solution: using segments

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Input: $(x_1, y_1), (x_2, y_2), \dots, (x_n, y_n), x_1 < x_2 < \dots < x_n$
penalty parameter $C > 0$

Segmented Least Squares

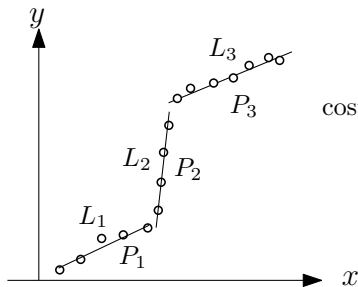
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minimize cost := error + penalty
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$$\begin{aligned} \text{cost} &= \text{error}(L_1, P_1) \\ &+ \text{error}(L_2, P_2) \\ &+ \text{error}(L_3, P_3) \\ &+ 3C \end{aligned}$$

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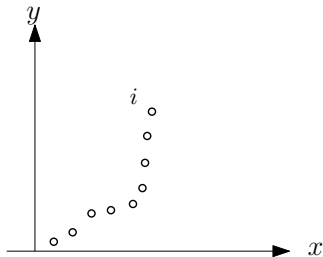
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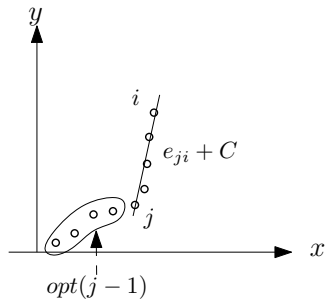
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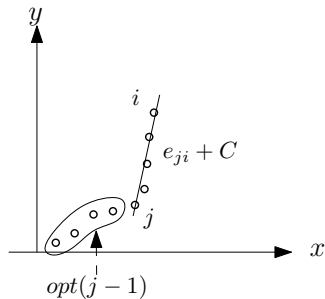
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- running time = $O(n^2)$.

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Input: an integer bound $W > 0$

a set of n items, each with an integer weight $w_i > 0$

Output: a subset S of items that

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- $W = 35, n = 5, w = (14, 9, 17, 10, 13)$

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Example:

- $W = 35, n = 5, w = (14, 9, 17, 10, 13)$
- Optimum: $S = \{1, 2, 4\}$ and $14 + 9 + 10 = 33$

Greedy Algorithms for Subset Sum

Candidate Algorithm:

- Sort according to non-increasing order of weights
- Select items in the order as long as the total weight remains below W

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- Consider the instance: $i, W', (w_1, w_2, \dots, w_i)$;
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Q: The value of the optimum solution that **does not contain** i ?

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A: $opt[i - 1, W']$

Q: The value of the optimum solution that **contains** i ?

A: $opt[i - 1, W' - w_i] + w_i$

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$$opt[i, W'] = \begin{cases} i = 0 \\ i > 0, w_i > W' \\ i > 0, w_i \leq W' \end{cases}$$

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Dynamic Programming

```
1: for  $W' \leftarrow 0$  to  $W$  do  
2:    $opt[0, W'] \leftarrow 0$   
3: for  $i \leftarrow 1$  to  $n$  do  
4:   for  $W' \leftarrow 0$  to  $W$  do  
5:      $opt[i, W'] \leftarrow opt[i - 1, W']$   
6:     if  $w_i \leq W'$  and  $opt[i - 1, W' - w_i] + w_i \geq opt[i, W']$  then  
7:        $opt[i, W'] \leftarrow opt[i - 1, W' - w_i] + w_i$   
8: return  $opt[n, W]$ 
```

Recover the Optimum Set

```
1: for  $W' \leftarrow 0$  to  $W$  do
2:    $opt[0, W'] \leftarrow 0$ 
3: for  $i \leftarrow 1$  to  $n$  do
4:   for  $W' \leftarrow 0$  to  $W$  do
5:      $opt[i, W'] \leftarrow opt[i - 1, W']$ 
6:      $b[i, W'] \leftarrow N$ 
7:     if  $w_i \leq W'$  and  $opt[i - 1, W' - w_i] + w_i \geq opt[i, W']$ 
   then
8:        $opt[i, W'] \leftarrow opt[i - 1, W' - w_i] + w_i$ 
9:        $b[i, W'] \leftarrow Y$ 
10: return  $opt[n, W]$ 
```

Recover the Optimum Set

```
1:  $i \leftarrow n, W' \leftarrow W, S \leftarrow \emptyset$   
2: while  $i > 0$  do  
3:   if  $b[i, W'] = Y$  then  
4:      $W' \leftarrow W' - w_i$   
5:      $S \leftarrow S \cup \{i\}$   
6:    $i \leftarrow i - 1$   
7: return  $S$ 
```

Running Time of Algorithm

```
1: for  $W' \leftarrow 0$  to  $W$  do  
2:    $opt[0, W'] \leftarrow 0$   
3: for  $i \leftarrow 1$  to  $n$  do  
4:   for  $W' \leftarrow 0$  to  $W$  do  
5:      $opt[i, W'] \leftarrow opt[i - 1, W']$   
6:     if  $w_i \leq W'$  and  $opt[i - 1, W' - w_i] + w_i \geq opt[i, W']$  then  
7:        $opt[i, W'] \leftarrow opt[i - 1, W' - w_i] + w_i$   
8: return  $opt[n, W]$ 
```

Running Time of Algorithm

```
1: for  $W' \leftarrow 0$  to  $W$  do
2:    $opt[0, W'] \leftarrow 0$ 
3: for  $i \leftarrow 1$  to  $n$  do
4:   for  $W' \leftarrow 0$  to  $W$  do
5:      $opt[i, W'] \leftarrow opt[i - 1, W']$ 
6:     if  $w_i \leq W'$  and  $opt[i - 1, W' - w_i] + w_i \geq opt[i, W']$  then
7:        $opt[i, W'] \leftarrow opt[i - 1, W' - w_i] + w_i$ 
8: return  $opt[n, W]$ 
```

- Running time is $O(nW)$

Running Time of Algorithm

```
1: for  $W' \leftarrow 0$  to  $W$  do
2:    $opt[0, W'] \leftarrow 0$ 
3: for  $i \leftarrow 1$  to  $n$  do
4:   for  $W' \leftarrow 0$  to  $W$  do
5:      $opt[i, W'] \leftarrow opt[i - 1, W']$ 
6:     if  $w_i \leq W'$  and  $opt[i - 1, W' - w_i] + w_i \geq opt[i, W']$  then
7:        $opt[i, W'] \leftarrow opt[i - 1, W' - w_i] + w_i$ 
8: return  $opt[n, W]$ 
```

- Running time is $O(nW)$
- Running time is **pseudo-polynomial** because it depends on value of the input integers.

Example

- $n = 4, w = (2, 3, 9, 8), W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1															
2															
3															
4															

Example

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0													
2															
3															
4															

Example

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2															
3															
4															

Example

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0													
3															
4															

Example

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0	2												
3															
4															

Example

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0	2	3	3										
3															
4															

Example

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0	2	3	3	5	5	5	5	5	5	5	5	5	5
3															
4															

Example

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0	2	3	3	5	5	5	5	5	5	5	5	5	5
3	0	0	2	3	3	5	5	5	5						
4															

Example

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0	2	3	3	5	5	5	5	5	5	5	5	5	5
3	0	0	2	3	3	5	5	5	5	9	9				
4															

Example

- $n = 4, w = (2, 3, 9, 8), W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0	2	3	3	5	5	5	5	5	5	5	5	5	5
3	0	0	2	3	3	5	5	5	5	9	9	11	12	12	14
4															

Example

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0	2	3	3	5	5	5	5	5	5	5	5	5	5
3	0	0	2	3	3	5	5	5	5	9	9	11	12	12	14
4	0	0	2	3	3	5	5	5							

Example

- $n = 4, w = (2, 3, 9, 8), W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0	2	3	3	5	5	5	5	5	5	5	5	5	5
3	0	0	2	3	3	5	5	5	5	9	9	11	12	12	14
4	0	0	2	3	3	5	5	5	8						

Example

- $n = 4, w = (2, 3, 9, 8), W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0	2	3	3	5	5	5	5	5	5	5	5	5	5
3	0	0	2	3	3	5	5	5	5	9	9	11	12	12	14
4	0	0	2	3	3	5	5	5	8	9					

Example

- $n = 4, w = (2, 3, 9, 8), W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0	2	3	3	5	5	5	5	5	5	5	5	5	5
3	0	0	2	3	3	5	5	5	5	9	9	11	12	12	14
4	0	0	2	3	3	5	5	5	8	9	10				

Example

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0	2	3	3	5	5	5	5	5	5	5	5	5	5
3	0	0	2	3	3	5	5	5	5	9	9	11	12	12	14
4	0	0	2	3	3	5	5	5	8	9	10	11	12		

Example

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0	2	3	3	5	5	5	5	5	5	5	5	5	5
3	0	0	2	3	3	5	5	5	5	9	9	11	12	12	14
4	0	0	2	3	3	5	5	5	8	9	10	11	12	13	

Example

- $n = 4, w = (2, 3, 9, 8), W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2
2	0	0	2	3	3	5	5	5	5	5	5	5	5	5	5
3	0	0	2	3	3	5	5	5	5	9	9	11	12	12	14
4	0	0	2	3	3	5	5	5	8	9	10	11	12	13	14

Avoiding Unnecessary Computation and Memory Using Memoized Algorithm and Hash Map

compute-opt(i, W')

```
1: if  $opt[i, W'] \neq \perp$  then return  $opt[i, W']$ 
2: if  $i = 0$  then  $r \leftarrow 0$ 
3: else
4:    $r \leftarrow \text{compute-opt}(i - 1, W')$ 
5:   if  $w_i \leq W'$  then
6:      $r' \leftarrow \text{compute-opt}(i - 1, W' - w_i) + w_i$ 
7:     if  $r' > r$  then  $r \leftarrow r'$ 
8:  $opt[i, W'] \leftarrow r$ 
9: return  $r$ 
```

- Use hash map for opt

Example Using Memoized Rounding

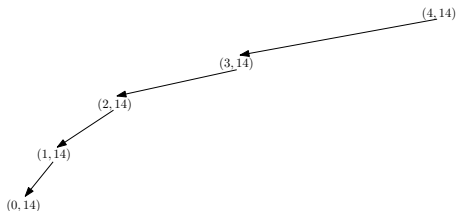
- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0															
1															
2															
3															
4															

Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

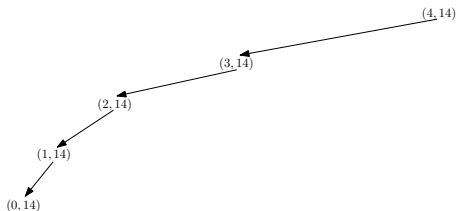
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0															
1															
2															
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

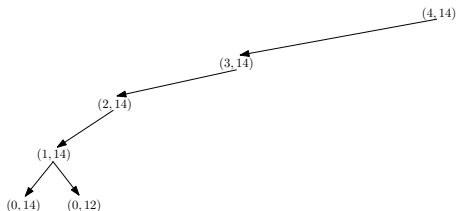
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0															0
1															
2															
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

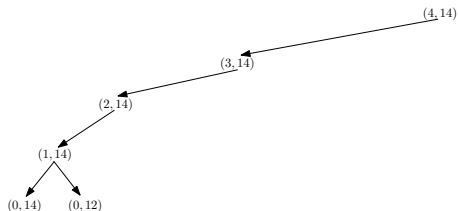
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0															0
1															
2															
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

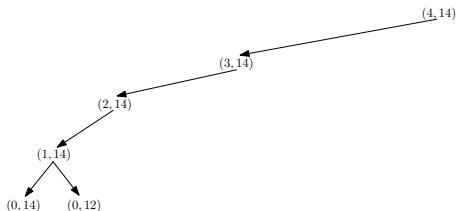
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0													0		0
1															
2															
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

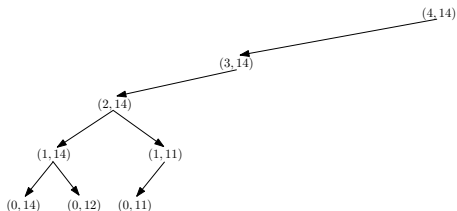
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0													0		0
1															2
2															
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

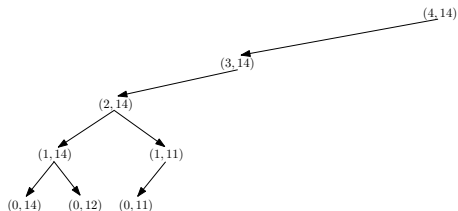
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0													0		0
1															2
2															
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

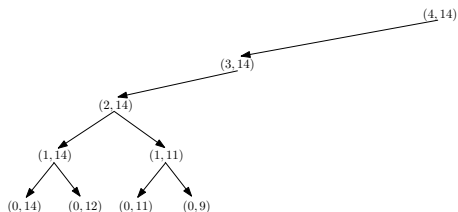
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0												0	0		0
1															2
2															
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

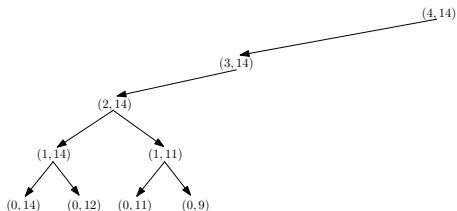
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0												0	0		0
1															2
2															
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

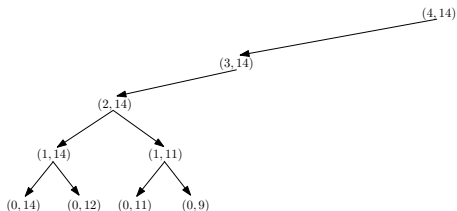
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0										0		0	0		0
1															2
2															
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

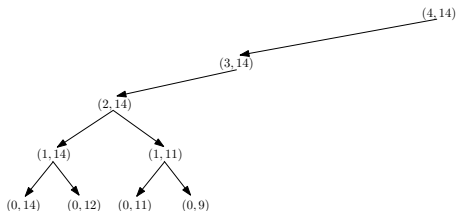
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0										0		0	0		0
1												2			2
2															
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

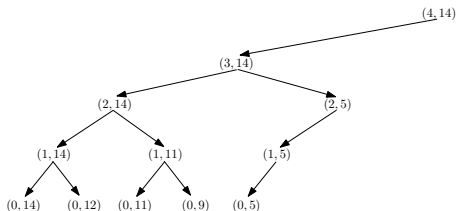
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0										0		0	0		0
1												2			2
2															5
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

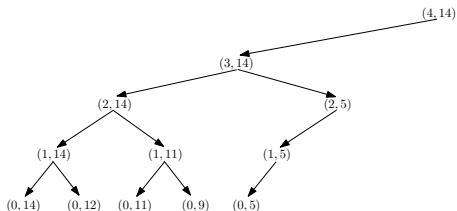
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0										0		0	0		0
1												2			2
2															5
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

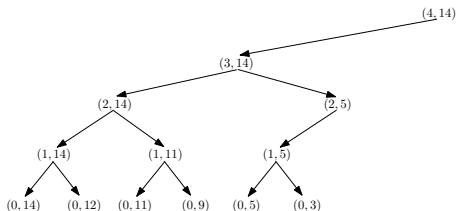
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0						0				0		0	0		0
1												2			2
2															5
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

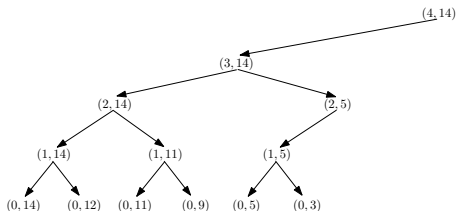
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0						0				0		0	0		0
1												2			2
2															5
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

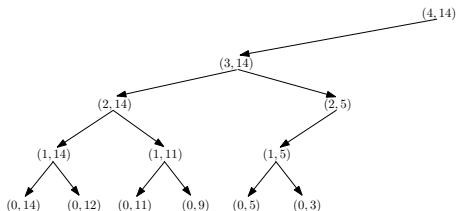
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0				0		0				0		0	0		0
1												2			2
2															5
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

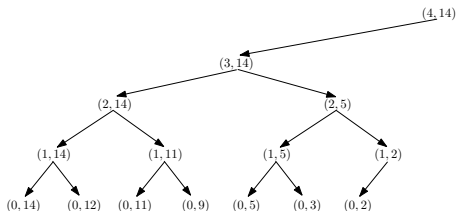
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0				0		0				0		0	0		0
1						2						2			2
2															5
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

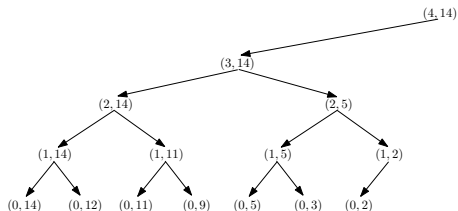
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0				0		0				0		0	0		0
1						2						2			2
2															5
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

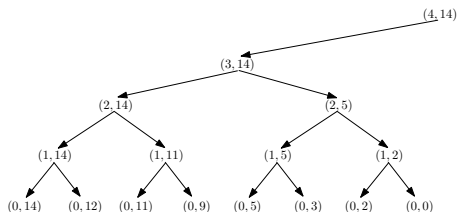
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0			0	0		0				0		0	0		0
1						2						2			2
2															5
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

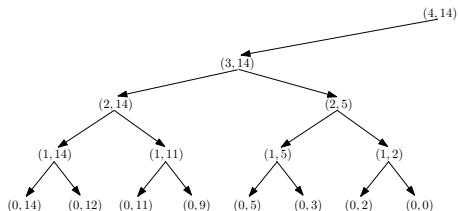
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0			0	0		0				0		0	0		0
1						2						2			2
2															5
3															
4															



Example Using Memoized Rounding

- $n = 4$, $w = (2, 3, 9, 8)$, $W = 14$

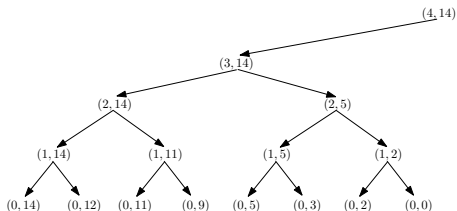
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0		0	0		0				0		0	0		0
1						2						2			2
2															5
3															
4															



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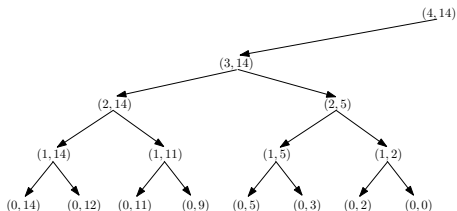
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0		0	0		0				0		0	0		0
1			2			2						2			2
2															5
3															
4															



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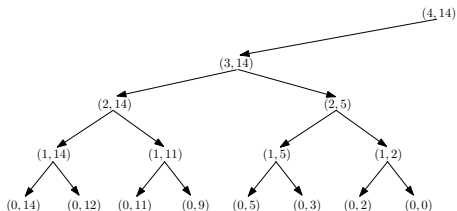
i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0		0	0		0				0		0	0		0
1			2			2						2			2
2						5									5
3															
4															



Example Using Memoized Rounding

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i, W'	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
0	0		0	0		0				0		0	0		0
1			2			2						2			2
2						5									5
3															14
4															



Outline

- 1 Weighted Interval Scheduling
- 2 Segmented Least Squares
- 3 Subset Sum Problem**
 - Related Problem: Knapsack Problem
- 4 Longest Common Subsequence
 - Longest Common Subsequence in Linear Space
- 5 Shortest Paths in Directed Acyclic Graphs
- 6 Matrix Chain Multiplication
- 7 Optimum Binary Search Tree
- 8 Summary

Knapsack Problem

Input: an integer bound $W > 0$

a set of n items, each with an integer weight $w_i > 0$

a value $v_i > 0$ for each item i

Output: a subset S of items that

$$\text{maximizes } \sum_{i \in S} v_i \quad \text{s.t. } \sum_{i \in S} w_i \leq W.$$

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- Motivation: you have budget W , and want to buy a subset of items of maximum total value

DP for Knapsack Problem

- $opt[i, W']$: the optimum value when budget is W' and items are $\{1, 2, 3, \dots, i\}$.
- If $i = 0$, $opt[i, W'] = 0$ for every $W' = 0, 1, 2, \dots, W$.

$$opt[i, W'] = \begin{cases} & i = 0 \\ & i > 0, w_i > W' \\ & i > 0, w_i \leq W' \end{cases}$$

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$$opt[i, W'] = \begin{cases} 0 & i = 0 \\ opt[i - 1, W'] & i > 0, w_i > W' \\ \max \left\{ \begin{array}{l} opt[i - 1, W'] \\ opt[i - 1, W' - w_i] + v_i \end{array} \right\} & i > 0, w_i \leq W' \end{cases}$$

Exercise: Items with 3 Parameters

Input: integer bounds $W > 0$, $Z > 0$,
a set of n items, each with an integer weight $w_i > 0$
a size $z_i > 0$ for each item i
a value $v_i > 0$ for each item i

Output: a subset S of items that

$$\begin{aligned} & \text{maximizes } \sum_{i \in S} v_i && \text{s.t.} \\ & \sum_{i \in S} w_i \leq W \text{ and } \sum_{i \in S} z_i \leq Z \end{aligned}$$

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Subsequence

- $A = bacdca$
- $C = adca$

Subsequence

- $A = bacdca$
- $C = adca$
- C is a subsequence of A

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Def. Given two sequences $A[1 .. n]$ and $C[1 .. t]$ of letters, C is called a **subsequence** of A if there exists integers $1 \leq i_1 < i_2 < i_3 < \dots < i_t \leq n$ such that $A[i_j] = C[j]$ for every $j = 1, 2, 3, \dots, t$.

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- Exercise: how to check if sequence C is a subsequence of A ?

Longest Common Subsequence

Input: $A[1 .. n]$ and $B[1 .. m]$

Output: the longest common subsequence of A and B

Example:

- $A = \text{'bacdca'}$
- $B = \text{'adbcda'}$

Longest Common Subsequence

Input: $A[1 .. n]$ and $B[1 .. m]$

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Example:

- $A = 'bacdca'$
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- $LCS(A, B) = 'adca'$

Longest Common Subsequence

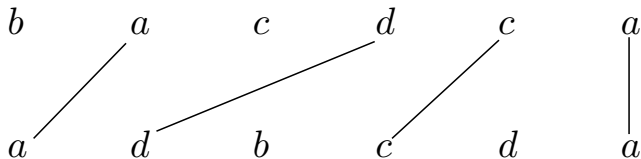
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Example:

- $A = 'bacdca'$
 - $B = 'adbcdca'$
 - $LCS(A, B) = 'adca'$
- Applications: edit distance (diff), similarity of DNAs

Matching View of LCS



- Goal of LCS: find a maximum-size non-crossing matching between letters in A and letters in B .

Reduce to Subproblems

- $A = \text{'bacdca'}$
- $B = \text{'adbcdca'}$

Reduce to Subproblems

- $A = \text{'bacdca'}$
- $B = \text{'adbcdca'}$

Reduce to Subproblems

- $A = \text{'bacdc'}$
- $B = \text{'adbcd'}$

Reduce to Subproblems

- $A = \text{'bacdc'}$
- $B = \text{'adbcd'}$
- either the last letter of A is not matched:
- or the last letter of B is not matched:

Reduce to Subproblems

- $A = 'bacdc'$
- $B = 'adbcd'$
- either the last letter of A is not matched:
 - need to compute $LCS('bacd', 'adbcd')$
- or the last letter of B is not matched:

Reduce to Subproblems

- $A = 'bacdc'$
- $B = 'adbcd'$
- either the last letter of A is not matched:
 - need to compute $LCS('bacd', 'adbcd')$
- or the last letter of B is not matched:
 - need to compute $LCS('bacdc', 'adbc')$

Dynamic Programming for LCS

- $opt[i, j], 0 \leq i \leq n, 0 \leq j \leq m$: length of longest common sub-sequence of $A[1 .. i]$ and $B[1 .. j]$.

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- if $i = 0$ or $j = 0$, then $opt[i, j] = 0$.
- if $i > 0, j > 0$, then

$$opt[i, j] = \begin{cases} & \text{if } A[i] = B[j] \\ & \text{if } A[i] \neq B[j] \end{cases}$$

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$$opt[i, j] = \begin{cases} opt[i - 1, j - 1] + 1 & \text{if } A[i] = B[j] \\ & \text{if } A[i] \neq B[j] \end{cases}$$

Dynamic Programming for LCS

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$$opt[i, j] = \begin{cases} opt[i - 1, j - 1] + 1 & \text{if } A[i] = B[j] \\ \max \begin{cases} opt[i - 1, j] \\ opt[i, j - 1] \end{cases} & \text{if } A[i] \neq B[j] \end{cases}$$

Dynamic Programming for LCS

```
1: for  $j \leftarrow 0$  to  $m$  do  
2:    $opt[0, j] \leftarrow 0$   
3: for  $i \leftarrow 1$  to  $n$  do  
4:    $opt[i, 0] \leftarrow 0$   
5:   for  $j \leftarrow 1$  to  $m$  do  
6:     if  $A[i] = B[j]$  then  
7:        $opt[i, j] \leftarrow opt[i - 1, j - 1] + 1$   
8:     else if  $opt[i, j - 1] \geq opt[i - 1, j]$  then  
9:        $opt[i, j] \leftarrow opt[i, j - 1]$   
10:    else  
11:       $opt[i, j] \leftarrow opt[i - 1, j]$ 
```

Dynamic Programming for LCS

```
1: for  $j \leftarrow 0$  to  $m$  do
2:    $opt[0, j] \leftarrow 0$ 
3: for  $i \leftarrow 1$  to  $n$  do
4:    $opt[i, 0] \leftarrow 0$ 
5:   for  $j \leftarrow 1$  to  $m$  do
6:     if  $A[i] = B[j]$  then
7:        $opt[i, j] \leftarrow opt[i - 1, j - 1] + 1, \pi[i, j] \leftarrow \text{“}\nearrow\text{”}$ 
8:     else if  $opt[i, j - 1] \geq opt[i - 1, j]$  then
9:        $opt[i, j] \leftarrow opt[i, j - 1], \pi[i, j] \leftarrow \text{“}\leftarrow\text{”}$ 
10:    else
11:       $opt[i, j] \leftarrow opt[i - 1, j], \pi[i, j] \leftarrow \text{“}\uparrow\text{”}$ 
```


Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥						
2	0 ⊥						
3	0 ⊥						
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						

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	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←					
2	0 ⊥						
3	0 ⊥						
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						

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0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←				
2	0 ⊥						
3	0 ⊥						
4	0 ⊥						
5	0 ⊥						
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0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖			
2	0 ⊥						
3	0 ⊥						
4	0 ⊥						
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6	0 ⊥						

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	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←		
2	0 ⊥						
3	0 ⊥						
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						

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	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	
2	0 ⊥						
3	0 ⊥						
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						

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0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥						
3	0 ⊥						
4	0 ⊥						
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0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥						
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	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖					
3	0 ⊥						
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						

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	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←				
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	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←			
3	0 ⊥						
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	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←		
3	0 ⊥						
4	0 ⊥						
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6	0 ⊥						

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	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	
3	0 ⊥						
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	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥						
4	0 ⊥						
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0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥						
4	0 ⊥						
5	0 ⊥						
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	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑					
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						

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	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←				
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						

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<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←			
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖		
4	0 ⊥						
5	0 ⊥						
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Example

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<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥						
5	0 ⊥						
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑					
5	0 ⊥						
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖				
5	0 ⊥						
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←			
5	0 ⊥						
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←		
5	0 ⊥						
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	
5	0 ⊥						
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥						
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥						
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑					
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑				
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←			
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖		
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥						

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖					

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑				

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←			

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑		

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	

Example

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

Example: Find Common Subsequence

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

Example: Find Common Subsequence

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

Example: Find Common Subsequence

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

Example: Find Common Subsequence

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

Example: Find Common Subsequence

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

Example: Find Common Subsequence

	1	2	3	4	5	6
A	b	a	c	d	c	a
B	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

Example: Find Common Subsequence

	1	2	3	4	5	6
A	b	a	c	d	c	a
B	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↖	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↖	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↖	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↖

Example: Find Common Subsequence

	1	2	3	4	5	6
<i>A</i>	b	a	c	d	c	a
<i>B</i>	a	d	b	c	d	a

	0	1	2	3	4	5	6
0	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥	0 ⊥
1	0 ⊥	0 ←	0 ←	1 ↖	1 ←	1 ←	1 ←
2	0 ⊥	1 ↘	1 ←	1 ←	1 ←	1 ←	2 ↖
3	0 ⊥	1 ↑	1 ←	1 ←	2 ↖	2 ←	2 ←
4	0 ⊥	1 ↑	2 ↘	2 ←	2 ←	3 ↖	3 ←
5	0 ⊥	1 ↑	2 ↑	2 ←	3 ↘	3 ←	3 ←
6	0 ⊥	1 ↖	2 ↑	2 ←	3 ↑	3 ←	4 ↘

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Find Common Subsequence

```
1:  $i \leftarrow n, j \leftarrow m, S \leftarrow ()$ 
2: while  $i > 0$  and  $j > 0$  do
3:   if  $\pi[i, j] = "\searrow"$  then
4:     add  $A[i]$  to beginning of  $S, i \leftarrow i - 1, j \leftarrow j - 1$ 
5:   else if  $\pi[i, j] = "\uparrow"$  then
6:      $i \leftarrow i - 1$ 
7:   else
8:      $j \leftarrow j - 1$ 
9: return  $S$ 
```

Variants of Problem

Edit Distance with Insertions and Deletions

Input: a string A

each time we can delete a letter from A or insert a letter to A

Output: minimum number of operations (insertions or deletions) we need to change A to B ?

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Obs. $\#OPs = \text{length}(A) + \text{length}(B) - 2 \cdot \text{length}(\text{LCS}(A, B))$

Variants of Problem

Edit Distance with Insertions, Deletions and Replacing

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Example:

- $A = \text{ocurrance}$, $B = \text{occurrence}$.
- 2 operations: insert 'c', change 'a' to 'e'
- Not related to LCS any more

Edit Distance with Replacing: Reduction to a Variant of LCS

- Need to match letters in A and B , every letter is matched at most once and there should be no crosses.
- However, we can **match two different letters**: Matching a same letter gives score 2, matching two different letters gives score 1.
- Need to maximize the score.
- DP recursion for the case $i > 0$ and $j > 0$:

$$opt[i, j] = \begin{cases} opt[i - 1, j - 1] + 2 & \text{if } A[i] = B[j] \\ \max \begin{cases} opt[i - 1, j] \\ opt[i, j - 1] \\ opt[i - 1, j - 1] + 1 \end{cases} & \text{if } A[i] \neq B[j] \end{cases}$$

- Relation : $\#OPs = \text{length}(A) + \text{length}(B) - \text{max_score}$

Edit Distance (with Replacing): using DP directly

- $opt[i, j], 0 \leq i \leq n, 0 \leq j \leq m$: edit distance between $A[1 .. i]$ and $B[1 .. j]$.

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Example:

- Input: acbcedeacab

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Example:

- Input: **acbc**ede**acab**
- Output: **aced**eca

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Computing the Length of LCS

```
1: for  $j \leftarrow 0$  to  $m$  do
2:    $opt[0, j] \leftarrow 0$ 
3: for  $i \leftarrow 1$  to  $n$  do
4:    $opt[i, 0] \leftarrow 0$ 
5:   for  $j \leftarrow 1$  to  $m$  do
6:     if  $A[i] = B[j]$  then
7:        $opt[i, j] \leftarrow opt[i - 1, j - 1] + 1$ 
8:     else if  $opt[i, j - 1] \geq opt[i - 1, j]$  then
9:        $opt[i, j] \leftarrow opt[i, j - 1]$ 
10:    else
11:       $opt[i, j] \leftarrow opt[i - 1, j]$ 
```

Obs. The i -th row of table only depends on $(i - 1)$ -th row.

Reducing Space to $O(n + m)$

Obs. The i -th row of table only depends on $(i - 1)$ -th row.

Q: How to use this observation to reduce space?

Reducing Space to $O(n + m)$

Obs. The i -th row of table only depends on $(i - 1)$ -th row.

Q: How to use this observation to reduce space?

A: We only keep two rows: the $(i - 1)$ -th row and the i -th row.

Linear Space Algorithm to Compute Length of LCS

```
1: for  $j \leftarrow 0$  to  $m$  do
2:    $opt[0, j] \leftarrow 0$ 
3: for  $i \leftarrow 1$  to  $n$  do
4:    $opt[i \bmod 2, 0] \leftarrow 0$ 
5:   for  $j \leftarrow 1$  to  $m$  do
6:     if  $A[i] = B[j]$  then
7:        $opt[i \bmod 2, j] \leftarrow opt[i - 1 \bmod 2, j - 1] + 1$ 
8:     else if  $opt[i \bmod 2, j - 1] \geq opt[i - 1 \bmod 2, j]$  then
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12: return  $opt[n \bmod 2, m]$ 
```

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- Only keep the last two rows: only know how to match $A[n]$

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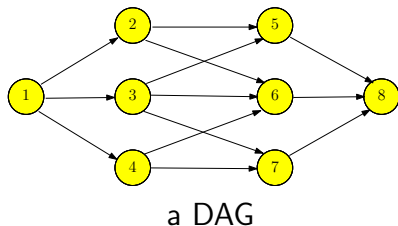
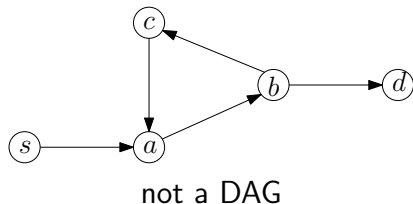
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- Using **Divide and Conquer** + Dynamic Programming:
 - Space: $O(m + n)$
 - Time: $O(nm)$

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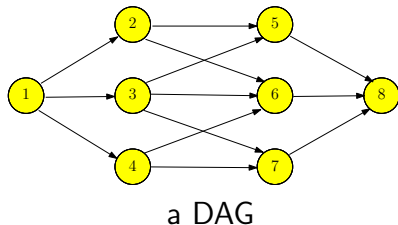
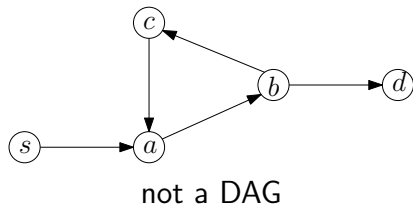
Directed Acyclic Graphs

Def. A directed acyclic graph (DAG) is a directed graph without (directed) cycles.



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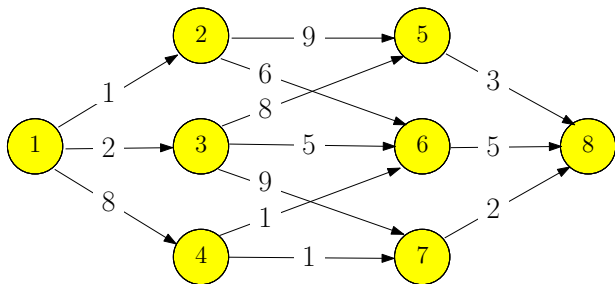
Lemma A directed graph is a DAG if and only if its vertices can be topologically sorted.

Shortest Paths in DAG

Input: directed acyclic graph $G = (V, E)$ and $w : E \rightarrow \mathbb{R}$.

Assume $V = \{1, 2, 3, \dots, n\}$ is topologically sorted: if $(i, j) \in E$, then $i < j$

Output: the shortest path from 1 to i , for every $i \in V$

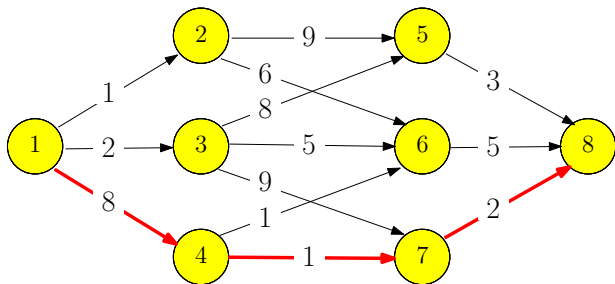


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Shortest Paths in DAG

- Use an adjacency list for incoming edges of each vertex i

Shortest Paths in DAG

```
1:  $f[1] \leftarrow 0$   
2: for  $i \leftarrow 2$  to  $n$  do  
3:    $f[i] \leftarrow \infty$   
4:   for each incoming edge  $(j, i)$  of  $i$  do  
5:     if  $f[j] + w(j, i) < f[i]$  then  
6:        $f[i] \leftarrow f[j] + w(j, i)$ 
```

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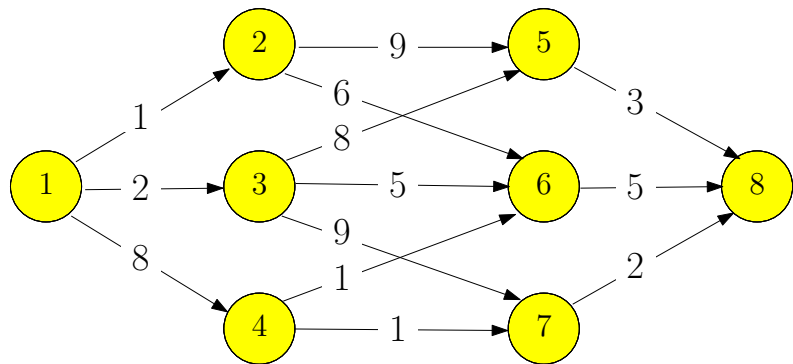
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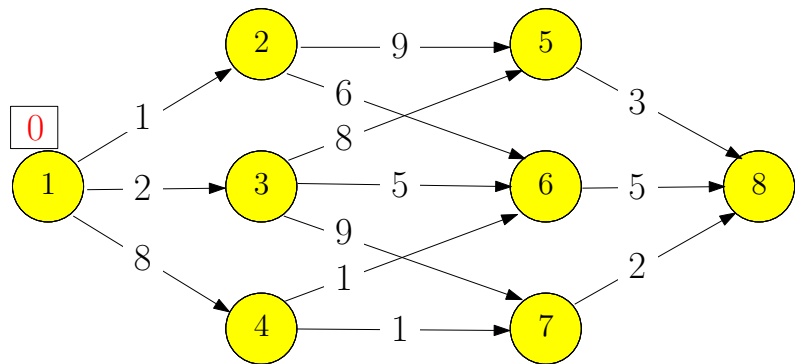
print-path(t)

```
1: if  $t = 1$  then
2:   print(1)
3:   return
4: print-path( $\pi(t)$ )
5: print(", ",  $t$ )
```

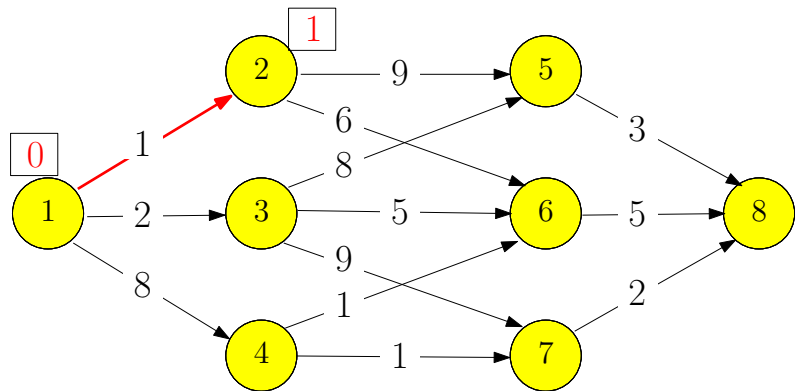
Example



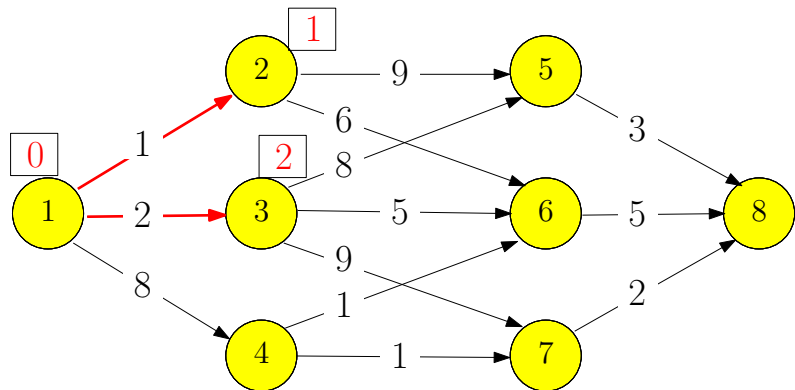
Example



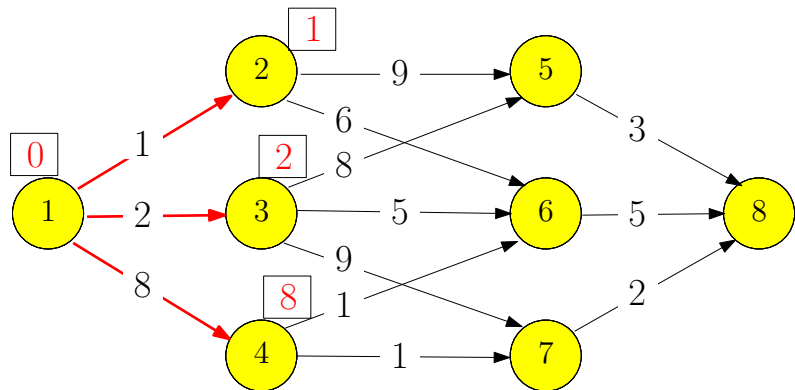
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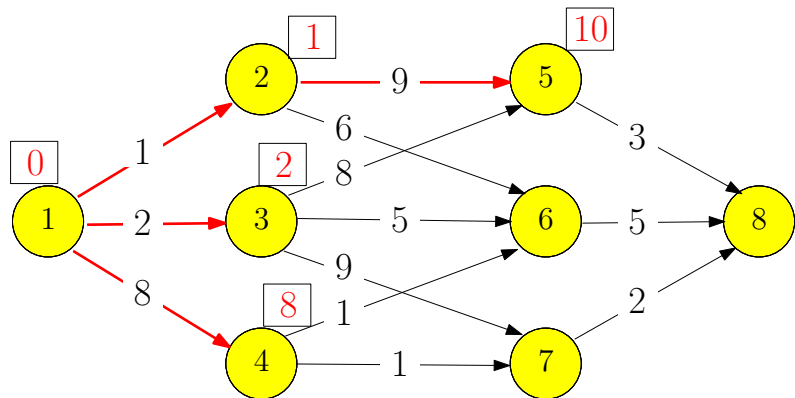
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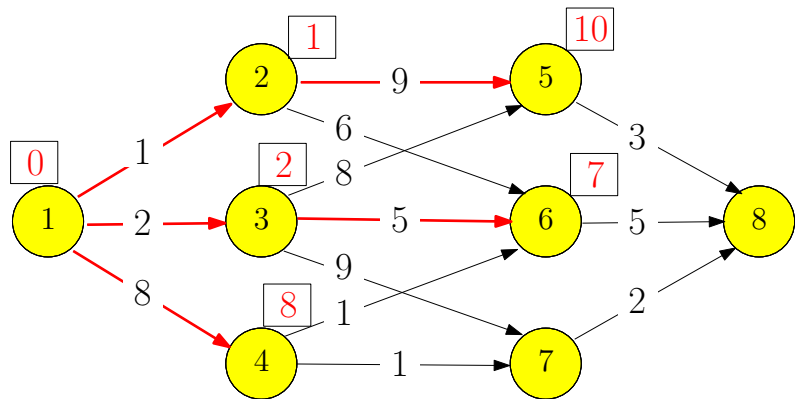
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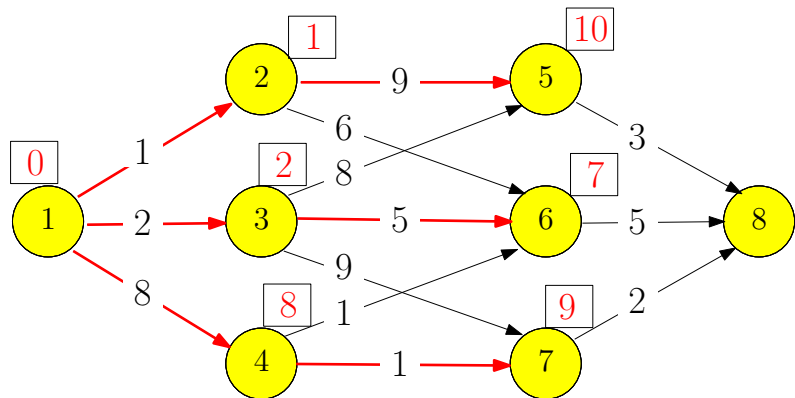
Example



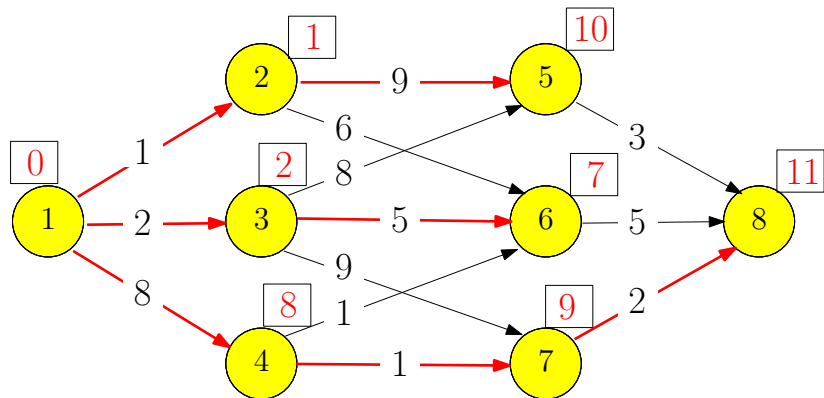
Example



Example



Example



Variant: Heaviest Path in a Directed Acyclic Graph

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Input: directed acyclic graph $G = (V, E)$ and $w : E \rightarrow \mathbb{R}$.

Assume $V = \{1, 2, 3, \dots, n\}$ is topologically sorted: if $(i, j) \in E$, then $i < j$

Output: the path with the **largest** weight (the **heaviest** path) from 1 to n .

- $f[i]$: weight of the **heaviest** path from 1 to i

$$f[i] = \begin{cases} & i = 1 \\ & i = 2, 3, \dots, n \end{cases}$$

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Matrix Chain Multiplication

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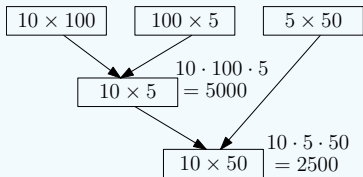
Input: n matrices A_1, A_2, \dots, A_n of sizes $r_1 \times c_1, r_2 \times c_2, \dots, r_n \times c_n$, such that $c_i = r_{i+1}$ for every $i = 1, 2, \dots, n - 1$.

Output: the order of computing $A_1 A_2 \dots A_n$ with the minimum number of multiplications

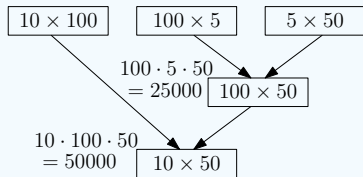
Fact Multiplying two matrices of size $r \times k$ and $k \times c$ takes $r \times k \times c$ multiplications.

Example:

- $A_1 : 10 \times 100$, $A_2 : 100 \times 5$, $A_3 : 5 \times 50$



$$\text{cost} = 5000 + 2500 = 7500$$

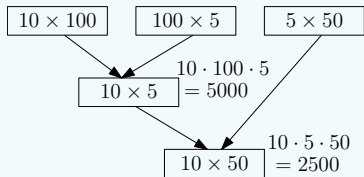


$$\text{cost} = 25000 + 50000 = 75000$$

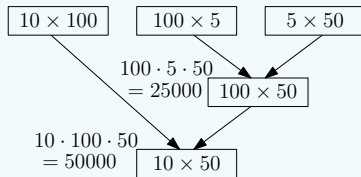
- $(A_1 A_2) A_3 : 10 \times 100 \times 5 + 10 \times 5 \times 50 = 7500$
- $A_1 (A_2 A_3) : 100 \times 5 \times 50 + 10 \times 100 \times 50 = 75000$

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- Assume the last step is $(A_1A_2 \cdots A_i)(A_{i+1}A_{i+2} \cdots A_n)$

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- Assume the last step is $(A_1A_2 \cdots A_i)(A_{i+1}A_{i+2} \cdots A_n)$
- Cost of last step: $r_1 \times c_i \times c_n$

Matrix Chain Multiplication: Design DP

- Assume the last step is $(A_1A_2 \cdots A_i)(A_{i+1}A_{i+2} \cdots A_n)$
- Cost of last step: $r_1 \times c_i \times c_n$
- Optimality for sub-instances: we need to compute $A_1A_2 \cdots A_i$ and $A_{i+1}A_{i+2} \cdots A_n$ optimally

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- $opt[i, j]$: the minimum cost of computing $A_iA_{i+1} \cdots A_j$

$$opt[i, j] = \begin{cases} 0 & i = j \\ \min_{k:i \leq k < j} (opt[i, k] + opt[k + 1, j] + r_i c_k c_j) & i < j \end{cases}$$

Matrix Chain Multiplication: Design DP

matrix-chain-multiplication($n, r[1..n], c[1..n]$)

- 1: let $opt[i, i] \leftarrow 0$ for every $i = 1, 2, \dots, n$
- 2: **for** $\ell \leftarrow 2$ to n **do**
- 3: **for** $i \leftarrow 1$ to $n - \ell + 1$ **do**
- 4: $j \leftarrow i + \ell - 1$
- 5: $opt[i, j] \leftarrow \infty$
- 6: **for** $k \leftarrow i$ to $j - 1$ **do**
- 7: **if** $opt[i, k] + opt[k + 1, j] + r_i c_k c_j < opt[i, j]$ **then**
- 8: $opt[i, j] \leftarrow opt[i, k] + opt[k + 1, j] + r_i c_k c_j$
- 9: **return** $opt[1, n]$

Recover the Optimum Way of Multiplication

matrix-chain-multiplication($n, r[1..n], c[1..n]$)

```
1: let  $opt[i, i] \leftarrow 0$  for every  $i = 1, 2, \dots, n$ 
2: for  $\ell \leftarrow 2$  to  $n$  do
3:   for  $i \leftarrow 1$  to  $n - \ell + 1$  do
4:      $j \leftarrow i + \ell - 1$ 
5:      $opt[i, j] \leftarrow \infty$ 
6:     for  $k \leftarrow i$  to  $j - 1$  do
7:       if  $opt[i, k] + opt[k + 1, j] + r_i c_k c_j < opt[i, j]$  then
8:          $opt[i, j] \leftarrow opt[i, k] + opt[k + 1, j] + r_i c_k c_j$ 
9:          $\pi[i, j] \leftarrow k$ 
10: return  $opt[1, n]$ 
```

Constructing Optimal Solution

Print-Optimal-Order(i, j)

```
1: if  $i = j$  then  
2:   print("A" $i$ )  
3: else  
4:   print("(")  
5:   Print-Optimal-Order( $i, \pi[i, j]$ )  
6:   Print-Optimal-Order( $\pi[i, j] + 1, j$ )  
7:   print(")")
```

matrix	A_1	A_2	A_3	A_4	A_5
size	3×5	5×2	2×6	6×9	9×4

$$\text{opt}[1, 2] = \text{opt}[1, 1] + \text{opt}[2, 2] + 3 \times 5 \times 2 = 30, \quad \pi[1, 2] = 1$$

$$\text{opt}[2, 3] = \text{opt}[2, 2] + \text{opt}[3, 3] + 5 \times 2 \times 6 = 60, \quad \pi[2, 3] = 2$$

$$\text{opt}[3, 4] = \text{opt}[3, 3] + \text{opt}[4, 4] + 2 \times 6 \times 9 = 108, \quad \pi[3, 4] = 3$$

$$\text{opt}[4, 5] = \text{opt}[4, 4] + \text{opt}[5, 5] + 6 \times 9 \times 4 = 216, \quad \pi[4, 5] = 4$$

$$\begin{aligned} \text{opt}[1, 3] &= \min\{\text{opt}[1, 1] + \text{opt}[2, 3] + 3 \times 5 \times 6, \\ &\quad \text{opt}[1, 2] + \text{opt}[3, 3] + 3 \times 2 \times 6\} \\ &= \min\{0 + 60 + 90, 30 + 0 + 36\} = 66, \quad \pi[1, 3] = 2 \end{aligned}$$

$$\begin{aligned} \text{opt}[2, 4] &= \min\{\text{opt}[2, 2] + \text{opt}[3, 4] + 5 \times 2 \times 9, \\ &\quad \text{opt}[2, 3] + \text{opt}[4, 4] + 5 \times 6 \times 9\} \\ &= \min\{0 + 108 + 90, 60 + 0 + 270\} = 198, \quad \pi[2, 4] = 2, \end{aligned}$$

matrix	A_1	A_2	A_3	A_4	A_5
size	3×5	5×2	2×6	6×9	9×4

$$\begin{aligned}
 \text{opt}[3, 5] &= \min\{\text{opt}[3, 3] + \text{opt}[4, 5] + 2 \times 6 \times 4, \\
 &\quad \text{opt}[3, 4] + \text{opt}[5, 5] + 2 \times 9 \times 4\} \\
 &= \min\{0 + 216 + 48, 108 + 0 + 72\} = 180,
 \end{aligned}$$

$$\pi[3, 5] = 4,$$

$$\begin{aligned}
 \text{opt}[1, 4] &= \min\{\text{opt}[1, 1] + \text{opt}[2, 4] + 3 \times 5 \times 9, \\
 &\quad \text{opt}[1, 2] + \text{opt}[3, 4] + 3 \times 2 \times 9, \\
 &\quad \text{opt}[1, 3] + \text{opt}[4, 4] + 3 \times 6 \times 9\} \\
 &= \min\{0 + 198 + 135, 30 + 108 + 54, 66 + 0 + 162\} = 192,
 \end{aligned}$$

$$\pi[1, 4] = 2,$$

matrix	A_1	A_2	A_3	A_4	A_5
size	3×5	5×2	2×6	6×9	9×4

$$\begin{aligned}
 \text{opt}[2, 5] &= \min\{\text{opt}[2, 2] + \text{opt}[3, 5] + 5 \times 2 \times 4, \\
 &\quad \text{opt}[2, 3] + \text{opt}[4, 5] + 5 \times 6 \times 4, \\
 &\quad \text{opt}[2, 4] + \text{opt}[5, 5] + 5 \times 9 \times 4\} \\
 &= \min\{0 + 180 + 40, 60 + 216 + 120, 198 + 0 + 180\} = 220,
 \end{aligned}$$

$$\begin{aligned}
 \text{opt}[1, 5] &= \min\{\text{opt}[1, 1] + \text{opt}[2, 5] + 3 \times 5 \times 4, \\
 &\quad \text{opt}[1, 2] + \text{opt}[3, 5] + 3 \times 2 \times 4, \\
 &\quad \text{opt}[1, 3] + \text{opt}[4, 5] + 3 \times 6 \times 4, \\
 &\quad \text{opt}[1, 4] + \text{opt}[5, 5] + 3 \times 9 \times 4\} \\
 &= \min\{0 + 220 + 60, 30 + 180 + 24, \\
 &\quad 66 + 216 + 72, 192 + 0 + 108\} \\
 &= 234,
 \end{aligned}$$

$$\pi[1, 5] = 2.$$

matrix	A_1	A_2	A_3	A_4	A_5
size	3×5	5×2	2×6	6×9	9×4

opt, π	$j = 1$	$j = 2$	$j = 3$	$j = 4$	$j = 5$
$i = 1$	0, /	30, 1	66, 2	192, 2	234, 2
$i = 2$		0, /	60, 2	198, 2	220, 2
$i = 3$			0, /	108, 3	180, 4
$i = 4$				0, /	216, 4
$i = 5$					0, /

opt, π	$j = 1$	$j = 2$	$j = 3$	$j = 4$	$j = 5$
$i = 1$	0, /	30, 1	66, 2	192, 2	234, 2
$i = 2$		0, /	60, 2	198, 2	220, 2
$i = 3$			0, /	108, 3	180, 4
$i = 4$				0, /	216, 4
$i = 5$					0, /

Print-Optimal-Order(1,5)

 Print-Optimal-Order(1, 2)

 Print-Optimal-Order(1, 1)

 Print-Optimal-Order(2, 2)

 Print-Optimal-Order(3, 5)

 Print-Optimal-Order(3, 4)

 Print-Optimal-Order(3, 3)

 Print-Optimal-Order(4, 4)

 Print-Optimal-Order(5, 5)

Optimum way for multiplication: $((A_1A_2)((A_3A_4)A_5))$

Outline

- 1 Weighted Interval Scheduling
- 2 Segmented Least Squares
- 3 Subset Sum Problem
 - Related Problem: Knapsack Problem
- 4 Longest Common Subsequence
 - Longest Common Subsequence in Linear Space
- 5 Shortest Paths in Directed Acyclic Graphs
- 6 Matrix Chain Multiplication
- 7 Optimum Binary Search Tree**
- 8 Summary

Optimum Binary Search Tree

- n elements $e_1 < e_2 < e_3 < \dots < e_n$
- e_i has frequency f_i
- goal: build a binary search tree for $\{e_1, e_2, \dots, e_n\}$ with the minimum accessing cost:

$$\sum_{i=1}^n f_i \times (\text{depth of } e_i \text{ in the tree})$$

Optimum Binary Search Tree

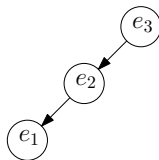
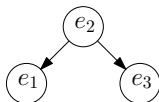
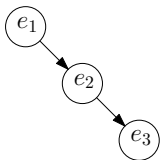
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- e_i has frequency f_i
- goal: build a binary search tree for $\{e_1, e_2, \dots, e_n\}$ with the minimum accessing cost:

$$\sum_{i=1}^n f_i \times (\text{depth of } e_i \text{ in the tree})$$

- motivation: the time to access e_i in the tree is linear in the depth of e_i

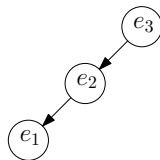
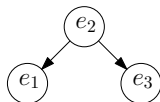
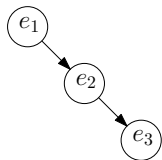
Optimum Binary Search Tree

- Example: $f_1 = 10, f_2 = 5, f_3 = 3$



Optimum Binary Search Tree

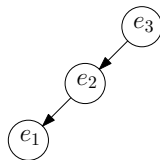
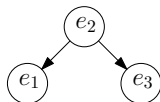
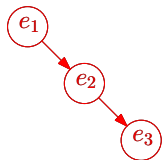
- Example: $f_1 = 10, f_2 = 5, f_3 = 3$



- $10 \times 1 + 5 \times 2 + 3 \times 3 = 29$
- $10 \times 2 + 5 \times 1 + 3 \times 2 = 31$
- $10 \times 3 + 5 \times 2 + 3 \times 1 = 43$

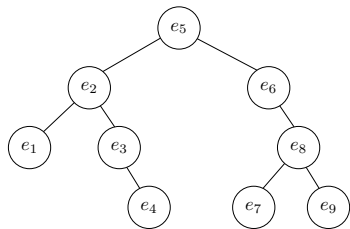
Optimum Binary Search Tree

- Example: $f_1 = 10, f_2 = 5, f_3 = 3$

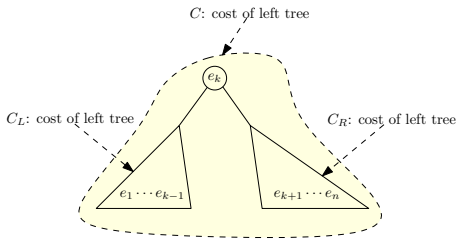


- $10 \times 1 + 5 \times 2 + 3 \times 3 = 29$
- $10 \times 2 + 5 \times 1 + 3 \times 2 = 31$
- $10 \times 3 + 5 \times 2 + 3 \times 1 = 43$

- suppose we decided to let e_k be the root
- e_1, e_2, \dots, e_{k-1} are on left sub-tree
- $e_{k+1}, e_{k+2}, \dots, e_n$ are on right sub-tree
- d_j : depth of e_j in our tree
- C, C_L, C_R : cost of tree, left sub-tree and right sub-tree



- $d_1 = 3, d_2 = 2, d_3 = 3, d_4 = 4, d_5 = 1,$
- $d_6 = 2, d_7 = 4, d_8 = 3, d_9 = 4,$
- $C = 3f_1 + 2f_2 + 3f_3 + 4f_4 + f_5 + 2f_6 + 4f_7 + 3f_8 + 4f_9$
- $C_L = 2f_1 + f_2 + 2f_3 + 3f_4$
- $C_R = f_6 + 3f_7 + 2f_8 + 3f_9$
- $C = C_L + C_R + \sum_{j=1}^9 f_j$



$$\begin{aligned}
 C &= \sum_{\ell=1}^n f_{\ell} d_{\ell} = \sum_{\ell=1}^n f_{\ell} (d_{\ell} - 1) + \sum_{\ell=1}^n f_{\ell} \\
 &= \sum_{\ell=1}^{k-1} f_{\ell} (d_{\ell} - 1) + \sum_{\ell=k+1}^n f_{\ell} (d_{\ell} - 1) + \sum_{\ell=1}^n f_{\ell} \\
 &= C_L + C_R + \sum_{\ell=1}^n f_{\ell}
 \end{aligned}$$

$$C = C_L + C_R + \sum_{\ell=1}^n f_{\ell}$$

- In order to minimize C , need to minimize C_L and C_R respectively

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- In order to minimize C , need to minimize C_L and C_R respectively
- $opt[i, j]$: the optimum cost for the instance $(f_i, f_{i+1}, \dots, f_j)$

$$opt[1, n] = (opt[1, k-1] + opt[k+1, n]) + \sum_{\ell=1}^n f_{\ell}$$

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- In order to minimize C , need to minimize C_L and C_R respectively
- $opt[i, j]$: the optimum cost for the instance $(f_i, f_{i+1}, \dots, f_j)$

$$opt[1, n] = \min_{k: 1 \leq k \leq n} (opt[1, k-1] + opt[k+1, n]) + \sum_{\ell=1}^n f_{\ell}$$

$$C = C_L + C_R + \sum_{\ell=1}^n f_{\ell}$$

- In order to minimize C , need to minimize C_L and C_R respectively
- $opt[i, j]$: the optimum cost for the instance $(f_i, f_{i+1}, \dots, f_j)$

$$opt[1, n] = \min_{k: 1 \leq k \leq n} (opt[1, k-1] + opt[k+1, n]) + \sum_{\ell=1}^n f_{\ell}$$

- In general, $opt[i, j] =$

$$\begin{cases} 0 & \text{if } i = j + 1 \\ \min_{k: i \leq k \leq j} (opt[i, k-1] + opt[k+1, j]) + \sum_{\ell=i}^j f_{\ell} & \text{if } i \leq j \end{cases}$$

Optimum Binary Search Tree

- 1: $fsum[0] \leftarrow 0$
- 2: **for** $i \leftarrow 1$ **to** n **do** $fsum[i] \leftarrow fsum[i - 1] + f_i$
 $\triangleright fsum[i] = \sum_{j=1}^i f_j$
- 3: **for** $i \leftarrow 0$ **to** n **do** $opt[i + 1, i] \leftarrow 0$
- 4: **for** $\ell \leftarrow 1$ **to** n **do**
- 5: **for** $i \leftarrow 1$ **to** $n - \ell + 1$ **do**
- 6: $j \leftarrow i + \ell - 1, opt[i, j] \leftarrow \infty$
- 7: **for** $k \leftarrow i$ **to** j **do**
- 8: **if** $opt[i, k - 1] + opt[k + 1, j] < opt[i, j]$ **then**
- 9: $opt[i, j] \leftarrow opt[i, k - 1] + opt[k + 1, j]$
- 10: $\pi[i, j] \leftarrow k$
- 11: $opt[i, j] \leftarrow opt[i, j] + fsum[j] - fsum[i - 1]$

Printing the Tree

Print-Tree(i, j)

```
1: if  $i > j$  then  
2:   return  
3: else  
4:   print('(')  
5:   Print-Tree( $i, \pi[i, j] - 1$ )  
6:   print( $\pi[i, j]$ )  
7:   Print-Tree( $\pi[i, j] + 1, j$ )  
8:   print('')
```


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Dynamic Programming

- Break up a problem into many **overlapping** sub-problems
- Build solutions for larger and larger sub-problems
- Use a **table** to store solutions for sub-problems for reuse

Comparison with greedy algorithms

- Greedy algorithm: each step is making a small progress towards constructing the solution
- Dynamic programming: the whole solution is constructed in the last step

Comparison with divide and conquer

- Divide and conquer: an instance is broken into many **independent** sub-instances, which are solved separately.
- Dynamic programming: the sub-instances we constructed are overlapping.

Definition of Cells for Problems We Learnt

- Weighted interval scheduling: $opt[i]$ = value of instance defined by jobs $\{1, 2, \dots, i\}$
- Segmented Least Square: $opt[i]$ = cost of instance defined by first i points.
- Subset sum, knapsack: $opt[i, W']$ = value of instance with items $\{1, 2, \dots, i\}$ and budget W'
- Longest common subsequence: $opt[i, j]$ = value of instance defined by $A[1..i]$ and $B[1..j]$
- Shortest paths in DAG: $f[v]$ = length of shortest path from s to v
- Matrix chain multiplication, optimum binary search tree:
 $opt[i, j]$ = value of instances defined by matrices i to j